

COMP 122/L Lecture 2

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Outline

- Operations on binary values
 - AND, OR, XOR, NOT
 - Bit shifting (left, two forms of right)
 - Addition
 - Subtraction
- Twos complement

Bitwise Operations

Bitwise AND

- Similar to logical AND (`& &`), except it works on a bit-by-bit manner
- Denoted by a single ampersand: `&`

```
(1001 &  
0101) =  
0001
```

Bitwise OR

- Similar to logical OR (`|`), except it works on a bit-by-bit manner
- Denoted by a single pipe character: `|`

```
(1001 |  
0101) =  
1101
```

Bitwise XOR

- Exclusive OR, denoted by a carat: \wedge
- Similar to bitwise OR, except that if both inputs are 1 then the result is 0

$$\begin{array}{r} (1001 \wedge \\ 0101) = \\ 1100 \end{array}$$

Bitwise NOT

- Similar to logical NOT (!), except it works on a bit-by-bit manner
- Denoted by a tilde character: ~

$$\begin{array}{r} \sim 1001 = \\ 0110 \end{array}$$

Shift Left

- Move all the bits N positions to the left, substituting in N 0s on the right

Shift Left

- Move all the bits N positions to the left, subbing in N 0s on the right

1001

Shift Left

- Move all the bits N positions to the left, subbing in N 0s on the right

$$\begin{array}{l} 1001 \ll 2 = \\ 100100 \end{array}$$

Shift Left

- Useful as a restricted form of multiplication
- Question: how?

1001 << 2 =
100100

Shift Left as Multiplication

- Equivalent decimal operation:

234

Shift Left as Multiplication

- Equivalent decimal operation:

$$\begin{array}{l} 234 \ll 1 = \\ 2340 \end{array}$$

Shift Left as Multiplication

- Equivalent decimal operation:

$$\begin{array}{l} 234 \ll 1 = \\ 2340 \end{array}$$

$$\begin{array}{l} 234 \ll 2 = \\ 23400 \end{array}$$

Multiplication

- Shifting left N positions multiplies by $(base)^N$
- Multiplying by 2 or 4 is often necessary (shift left 1 or 2 positions, respectively)
- Often a whoooole lot faster than telling the processor to multiply
- Compilers try hard to do this

$$\begin{array}{r} 234 \ll 2 = \\ 23400 \end{array}$$

Shift Right

- Move all the bits N positions to the right, subbing in **either** N 0s or N 1s on the left
- Two different forms

Shift Right

- Move all the bits N positions to the right, subbing in **either** N 0s or N (whatever the leftmost bit is)s on the left

- Two different forms

1001 >> 2 =

either 0010 **or** 1110

Shift Right Trick

- Question: If shifting left multiplies, what does shift right do?

Shift Right Trick

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 - Answer: divides in a similar way, but truncates result

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234

Shift Right Trick

- Question: If shifting left multiplies, what does shift right do?
 - Answer: divides in a similar way, but truncates result

$$234 \gg 1 = 23$$

Two Forms of Shift

Right

- Subbing in 0s makes sense
- What about subbing in the leftmost bit?
 - And why is this called “arithmetic” shift right?

1100 (arithmetic) >> 1 =
1110

Answer...Sort of

- Arithmetic form is intended for numbers in *twos complement*, whereas the non-arithmetic form is intended for *unsigned numbers*

Twos Complement

Problem

- Binary representation so far makes it easy to represent positive numbers and zero
- Question: What about representing negative numbers?

Twos Complement

- Way to represent positive integers, negative integers, and zero
- If 1 is in the *most significant bit* (generally leftmost bit in this class), then it is negative

Decimal to Twos Complement

- Example: -5 decimal to binary (twos complement)

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- First, convert the magnitude to an unsigned representation

Decimal to Twos Complement

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- First, convert the magnitude to an unsigned representation

$$5 \text{ (decimal)} = 0101 \text{ (binary)}$$

Decimal to Twos Complement

- Then, take the bits, and negate them

Decimal to Twos Complement

- Then, take the bits, and negate them

0101

Decimal to Twos Complement

- Then, take the bits, and negate them

$$\begin{array}{r} \sim 0101 = \\ 1010 \end{array}$$

Decimal to Twos Complement

- Finally, add one:

Decimal to Twos Complement

- Finally, add one:

1010

Decimal to Twos Complement

- Finally, add one:

$$\begin{array}{r} 1010 \\ + 1 \\ \hline 1011 \end{array} =$$

Twos Complement to Decimal

- Same operation: negate the bits, and add one

Twos Complement to Decimal

- Same operation: negate the bits, and add one

1011

Twos Complement to Decimal

- Same operation: negate the bits, and add one

$$\begin{array}{r} \sim 1011 = \\ 0100 \end{array}$$

Twos Complement to Decimal

- Same operation: negate the bits, and add one

0100

Twos Complement to Decimal

- Same operation: negate the bits, and add one

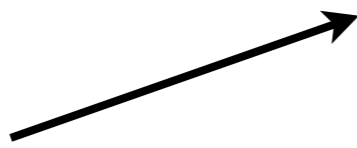
$$\begin{array}{r} 0100 + 1 = \\ 0101 \end{array}$$

Twos Complement to Decimal

- Same operation: negate the bits, and add one

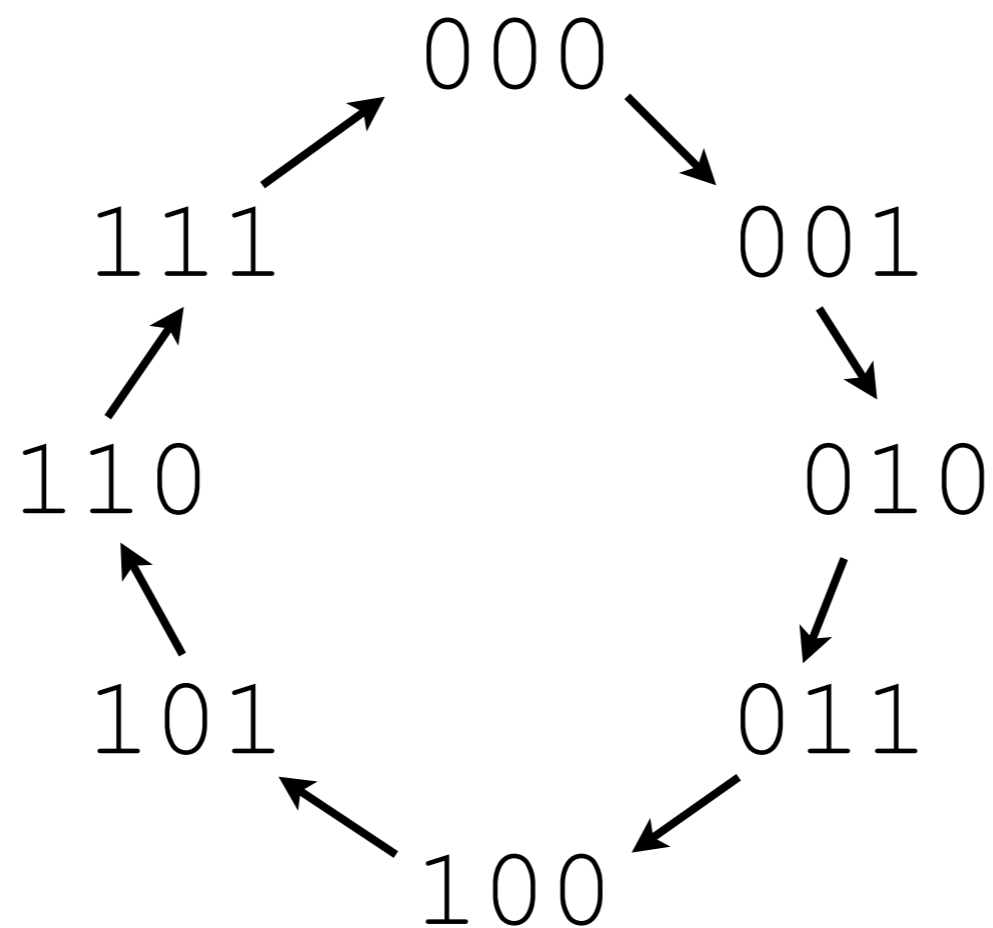
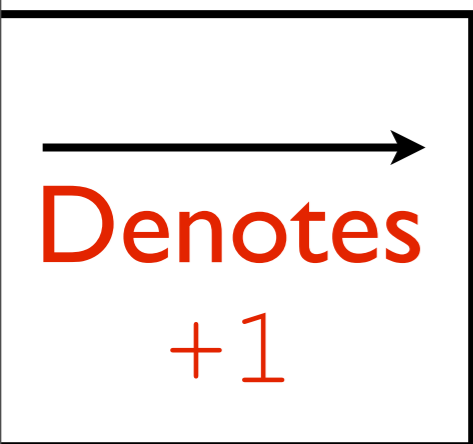
$$\begin{array}{r} 0100 + 1 = \\ 0101 = \\ -5 \end{array}$$

We started with
1011 - negative



Intuition

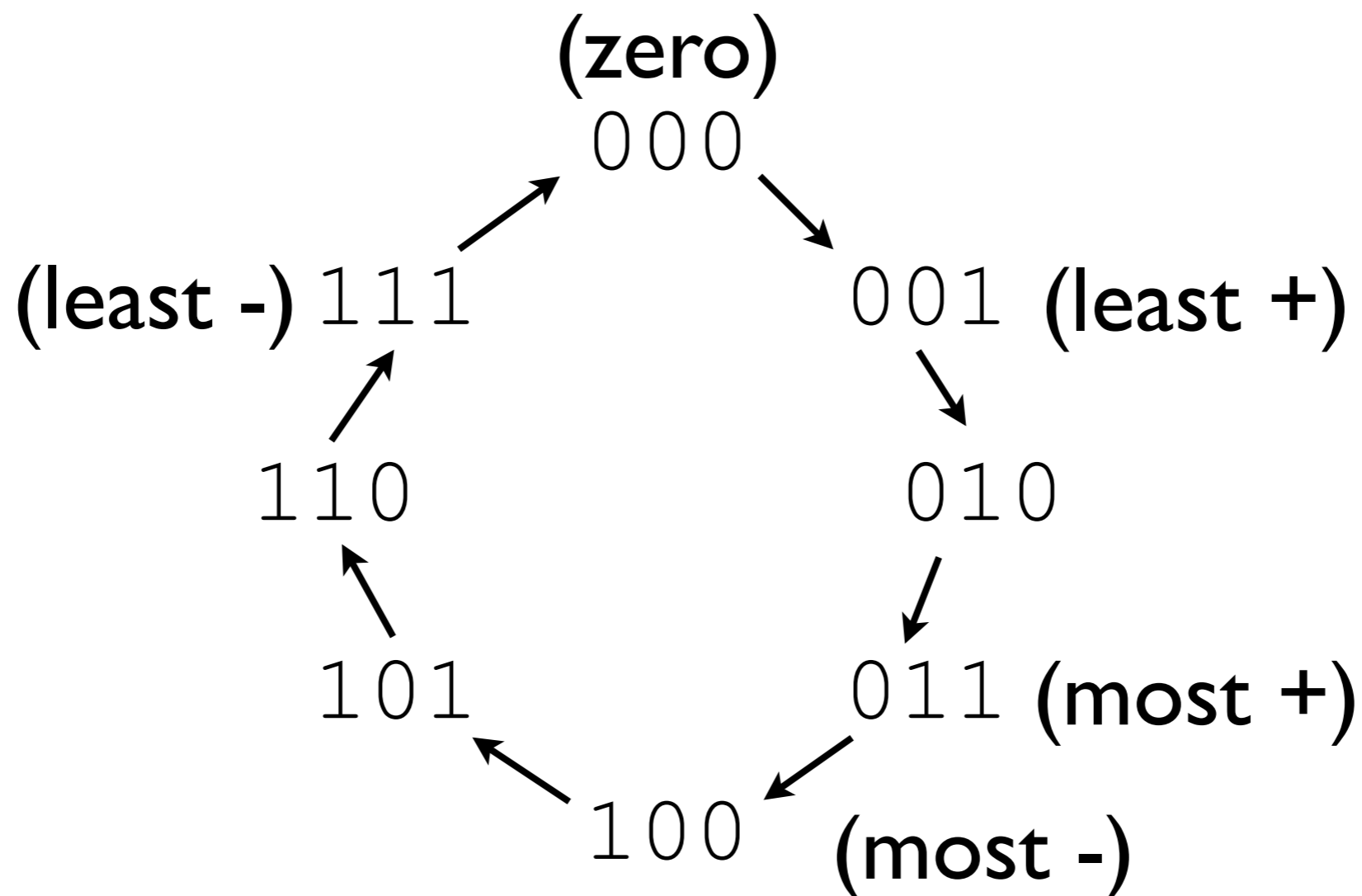
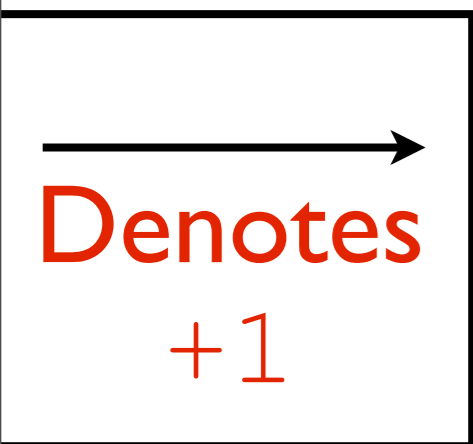
- Modular arithmetic, with the convention that a leading 1 bit means negative



-This is the intuition from Wikipedia, which makes a whole lot more sense

Intuition

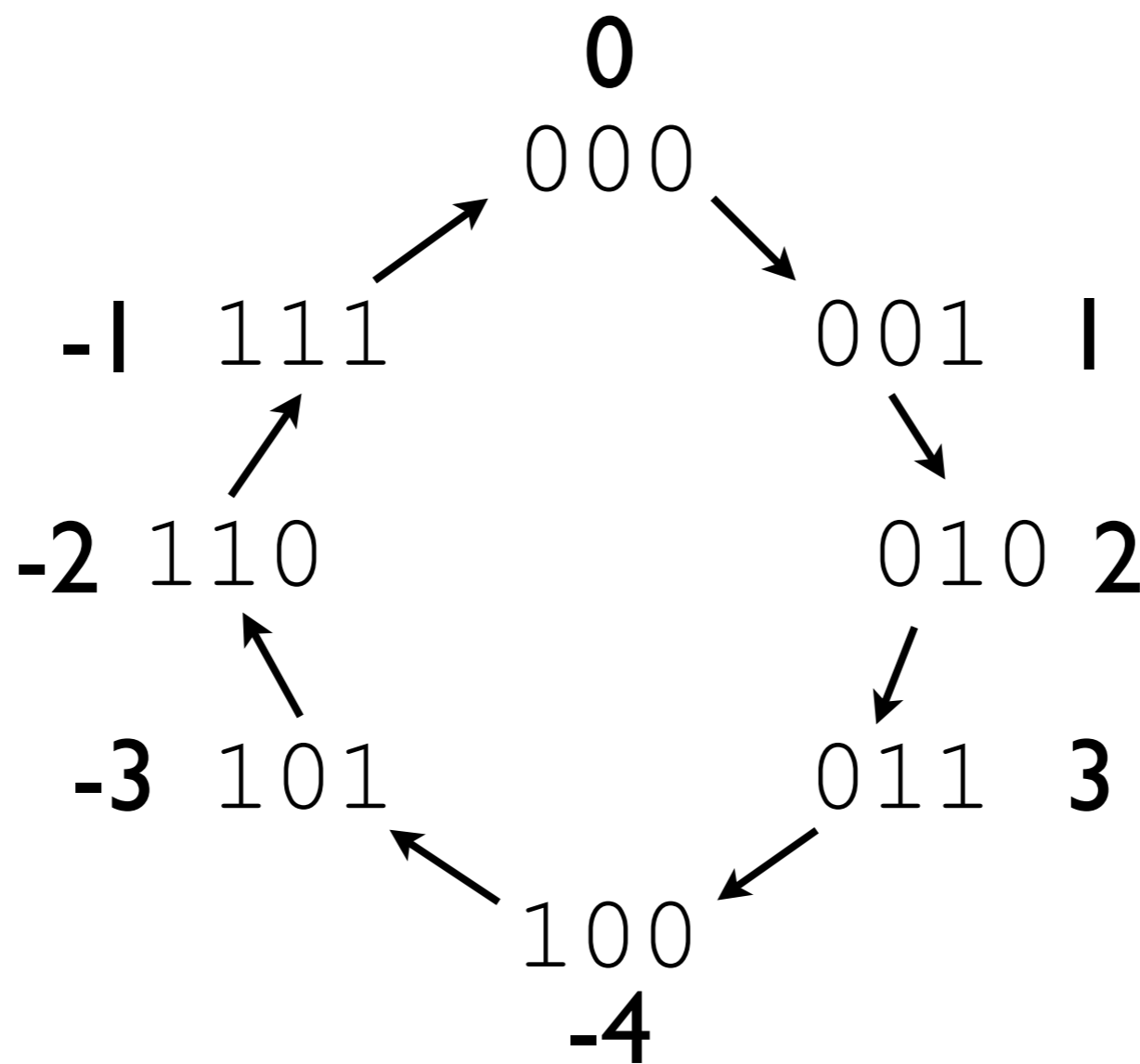
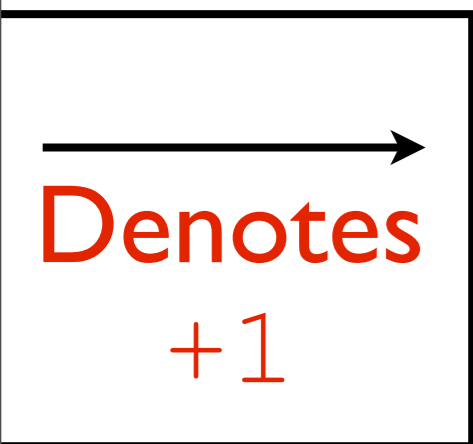
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- This is the intuition from Wikipedia, which makes a whole lot more sense
- There is still a lot of detail missing here – it's not necessary to understand in order to work with this. There is actually quite a bit of mathematics behind why this works

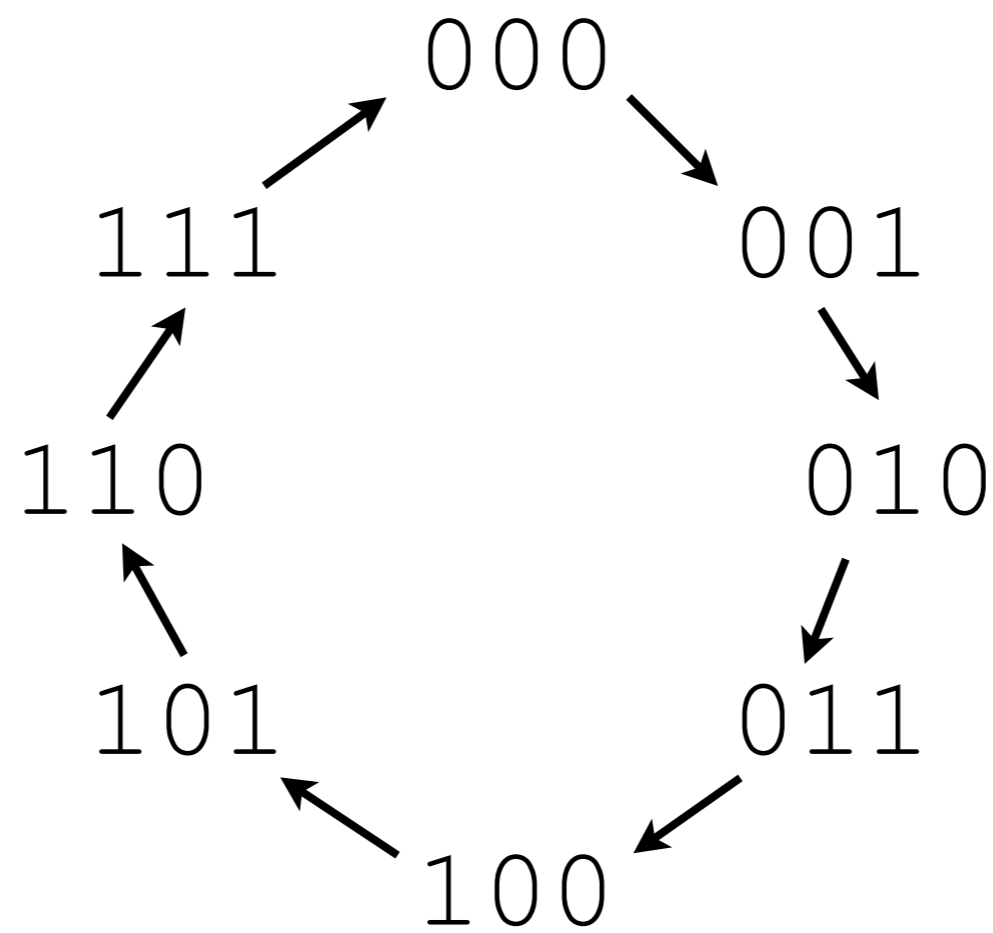
Intuition

- Modular arithmetic, with the convention that a leading 1 bit means negative



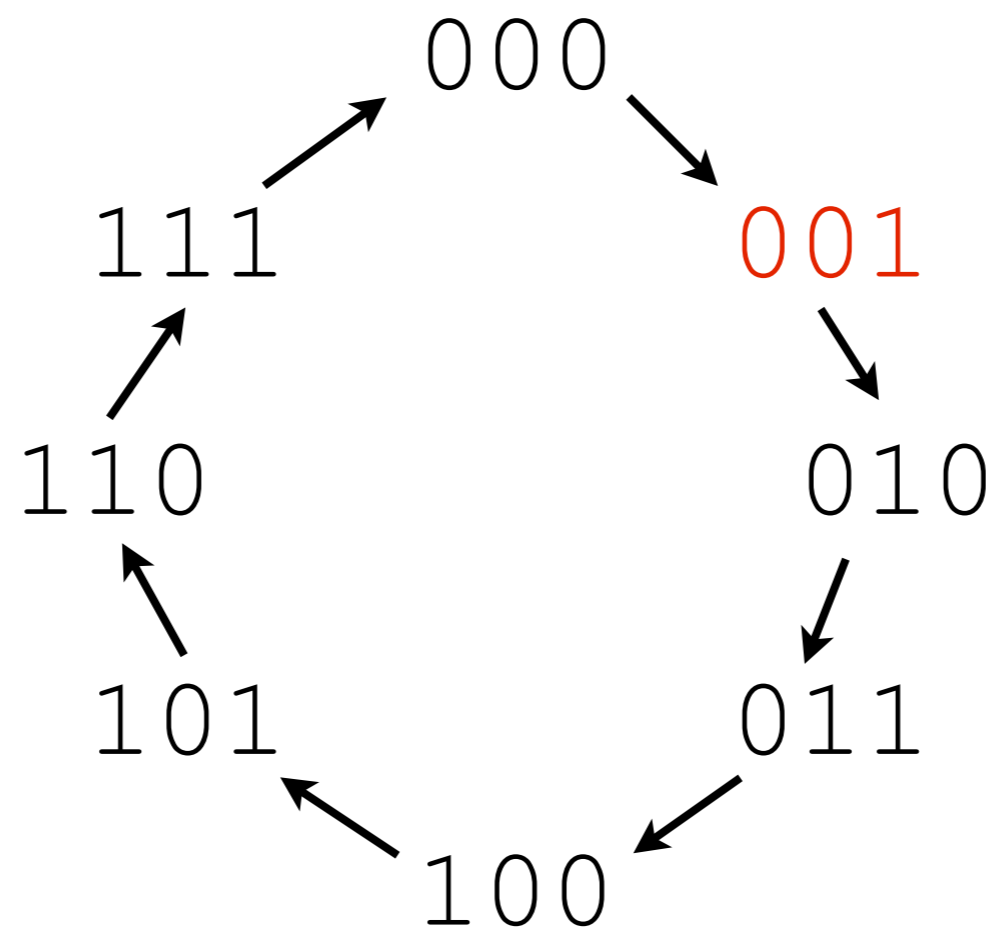
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Negation of 1



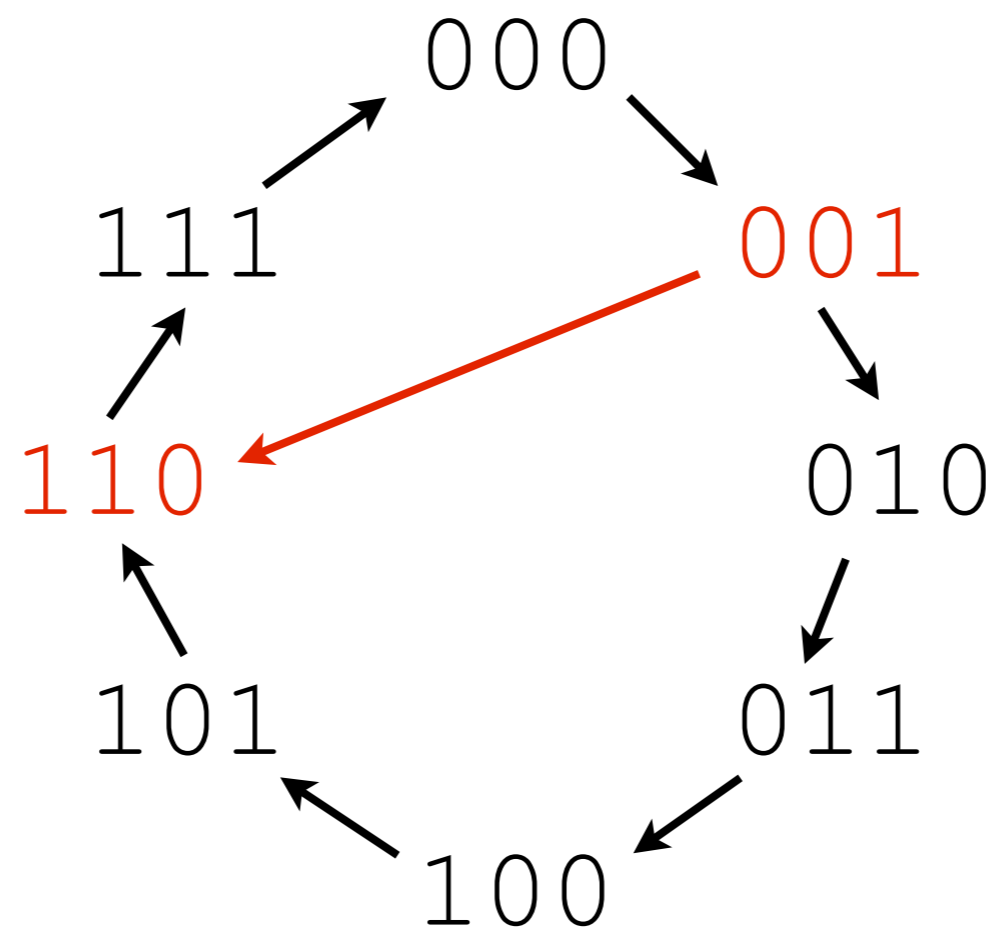
-Take our wheel from before

Negation of 1



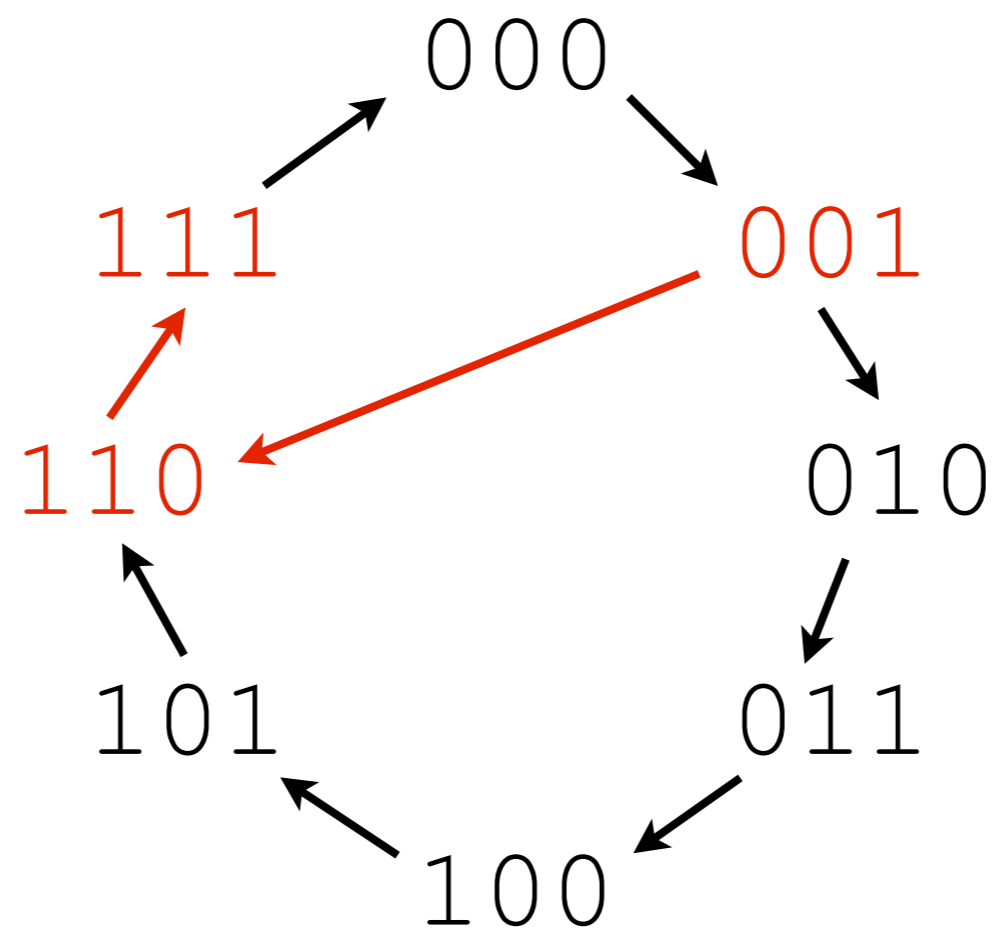
-This is 1

Negation of 1



-Inverted bits

Negation of 1

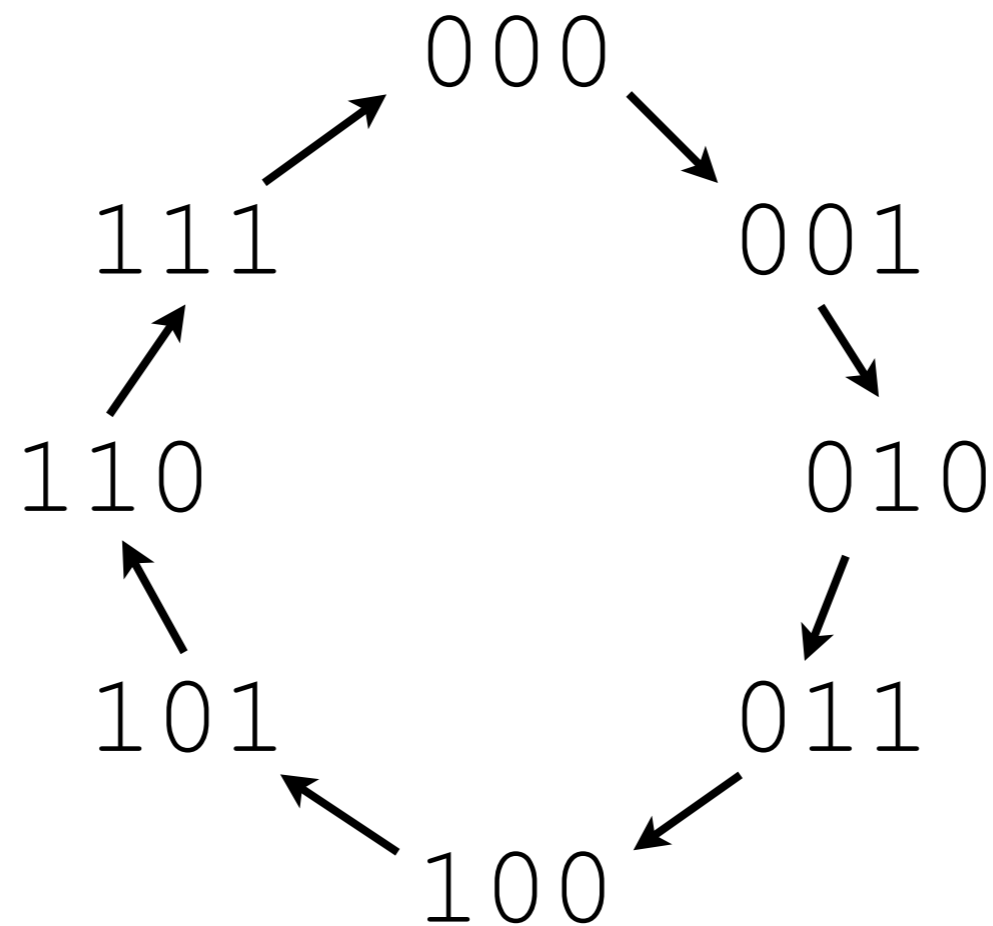


-Add 1

-This is exactly what we expected - binary 111 represents decimal -1

Consequences

- What is the negation of 000?



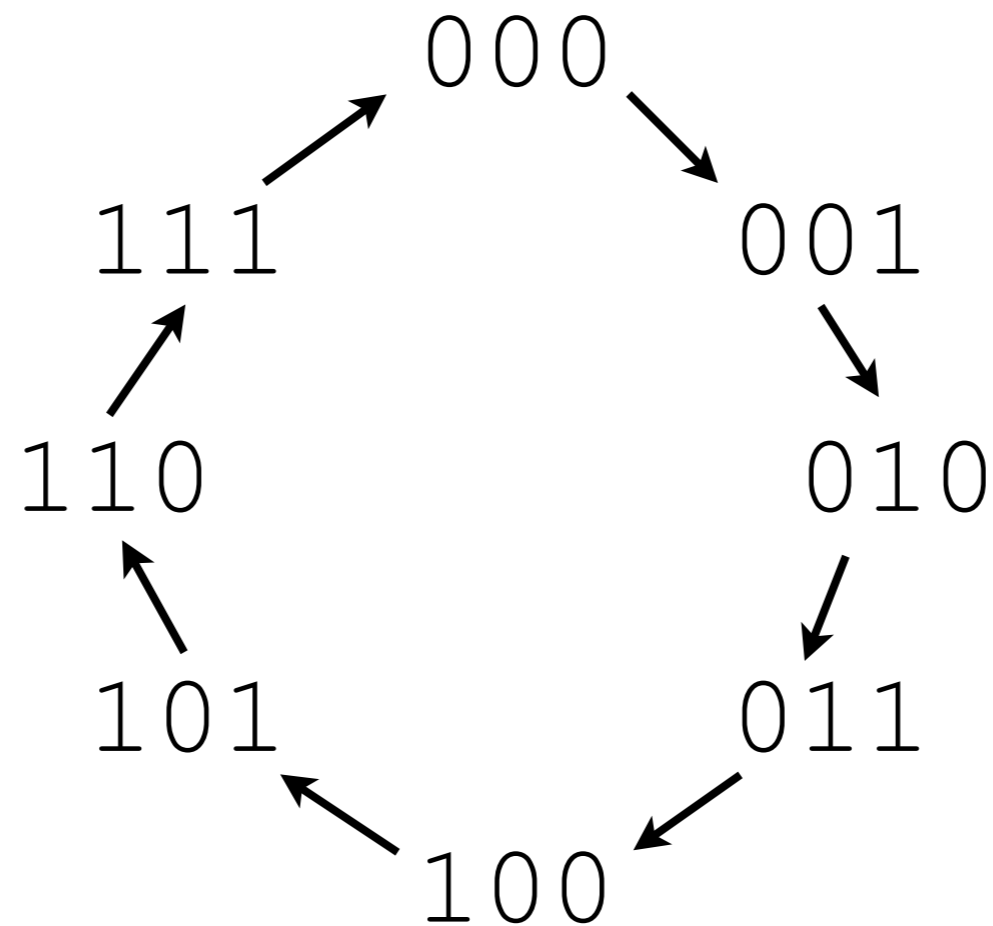
-Negate all bits: 000 -> 111

-Add one: 000

-Technically, adding one resulted in 1000, but that got cut off

Consequences

- What is the negation of 100?



-Negate all bits: 100 \rightarrow 011

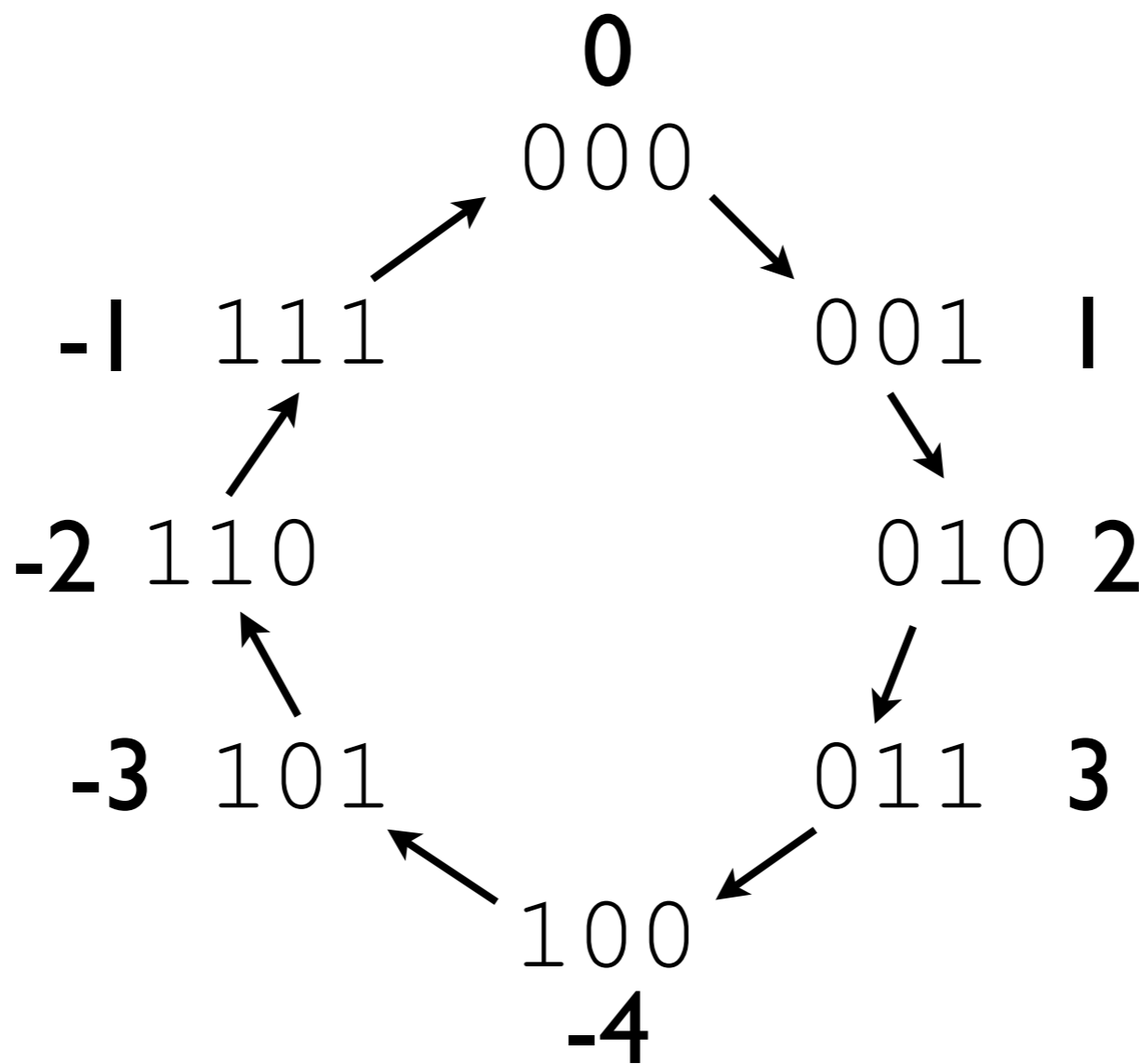
-Add one: 100

-Uh oh...this states that the negation of -4 is -4.

-Underlying problem is that we don't have a representation for 4 with just three bits

Arithmetic Shift Right

- **Not exactly** division by a power of two
- Consider $-3 / 2$



-101 (-3) shifted right yields 110 (-2), NOT 111 (-1) as expected from typical integer division

-Integer division rounds towards zero, whereas shift right rounds towards negative infinity

-This means they work identically for positive values, but not for negative values (also meaning they are always the same for unsigned values)

Addition

Building Up Addition

- Question: how might we add the following, in decimal?

$$\begin{array}{r} 986 \\ +123 \\ \hline \end{array}$$

?

Building Up Addition

- Question: how might we add the following, in decimal?

$$\begin{array}{r} 986 \\ +123 \\ \hline \end{array}$$

?

			$\begin{array}{r} 6 \\ +3 \\ \hline \end{array}$
--	--	--	--

Building Up Addition

- Question: how might we add the following, in decimal?

$$\begin{array}{r} 986 \\ +123 \\ \hline \end{array}$$

?

$$\begin{array}{r} 8 \\ +2 \\ \hline \end{array}$$

?

$$\begin{array}{r} 6 \\ +3 \\ \hline \end{array}$$

9

Building Up Addition

- Question: how might we add the following, in decimal?

$$\begin{array}{r} 986 \\ +123 \\ \hline \end{array}$$

?

Carry: 1

$$\begin{array}{r} 8 \\ +2 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 6 \\ +3 \\ \hline 9 \end{array}$$

Building Up Addition

- Question: how might we add the following, in decimal?

$$\begin{array}{r} 986 \\ +123 \\ \hline \end{array}$$

?

	$\begin{array}{r} 1 \\ 9 \\ +1 \\ \hline \end{array}$ <p style="text-align: center;">?</p>		
--	--	--	--

	$\begin{array}{r} 8 \\ +2 \\ \hline \end{array}$ <p style="text-align: center;">0</p>		
--	---	--	--

	$\begin{array}{r} 6 \\ +3 \\ \hline \end{array}$ <p style="text-align: center;">9</p>		
--	---	--	--

Building Up Addition

- Question: how might we add the following, in decimal?

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?

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Building Up Addition

- Question: how might we add the following, in decimal?

$$\begin{array}{r} 986 \\ +123 \\ \hline \end{array}$$

?

$$\begin{array}{r} 1 \\ +0 \\ \hline 1 \end{array}$$

$$\begin{array}{r} 1 \\ 9 \\ +1 \\ \hline 1 \end{array}$$

$$\begin{array}{r} 8 \\ +2 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 6 \\ +3 \\ \hline 9 \end{array}$$

Core Concepts

- We have a “primitive” notion of adding single digits, along with an idea of *carrying* digits
- We can build on this notion to add numbers together that are more than one digit long

Now in Binary

- Arguably simpler - fewer one-bit possibilities

0	0	1	1
+0	+1	+0	+1
--	--	--	--
?	?	?	?

Now in Binary

- Arguably simpler - fewer one-bit possibilities

0	0	1	1
+0	+1	+0	+1
--	--	--	--
0	1	1	0
			Carry: 1

Chaining the Carry

- Also need to account for any input carry

$\begin{array}{r} 0 \\ 0 \\ +0 \\ \hline 0 \end{array}$	$\begin{array}{r} 0 \\ 0 \\ +1 \\ \hline 1 \end{array}$	$\begin{array}{r} 0 \\ 1 \\ +0 \\ \hline 1 \end{array}$	$\begin{array}{r} 0 \\ 1 \\ +1 \\ \hline 0 \end{array} \text{ Carry: } 1$
$\begin{array}{r} 1 \\ 0 \\ +0 \\ \hline 1 \end{array}$	$\begin{array}{r} 1 \\ 0 \\ +1 \\ \hline 0 \end{array} \text{ Carry: } 1$	$\begin{array}{r} 1 \\ 1 \\ +0 \\ \hline 0 \end{array} \text{ Carry: } 1$	$\begin{array}{r} 1 \\ 1 \\ +1 \\ \hline 1 \end{array} \text{ Carry: } 1$

Adding Multiple Bits

- How might we add the numbers below?

```
  011
+001
-----
```


Adding Multiple Bits

- How might we add the numbers below?

$$\begin{array}{r} 0 \\ 011 \\ +001 \\ \hline \end{array}$$

-Need an initial carry-in of zero

Adding Multiple Bits

- How might we add the numbers below?

$$\begin{array}{r} 10 \\ 011 \\ +001 \\ \hline 0 \end{array}$$

-Need an initial carry-in of zero

Adding Multiple Bits

- How might we add the numbers below?

$$\begin{array}{r} 110 \\ 011 \\ +001 \\ \hline 00 \end{array}$$

-Need an initial carry-in of zero

Adding Multiple Bits

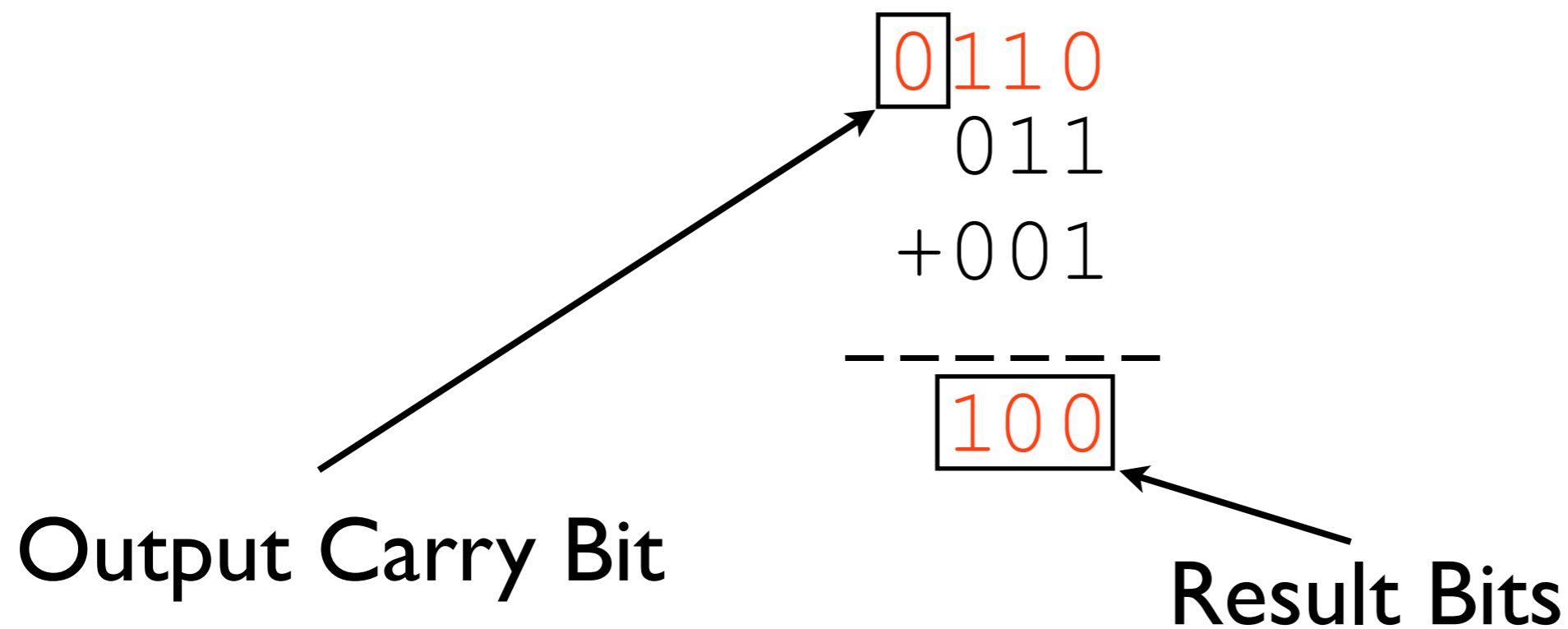
- How might we add the numbers below?

$$\begin{array}{r} 0110 \\ 011 \\ +001 \\ \hline 100 \end{array}$$

-Need an initial carry-in of zero

Adding Multiple Bits

- How might we add the numbers below?



-Need an initial carry-in of zero

Another Example

```
  111  
+001  
-----
```

Another Example

$$\begin{array}{r} 0 \\ 111 \\ +001 \\ \hline \end{array}$$

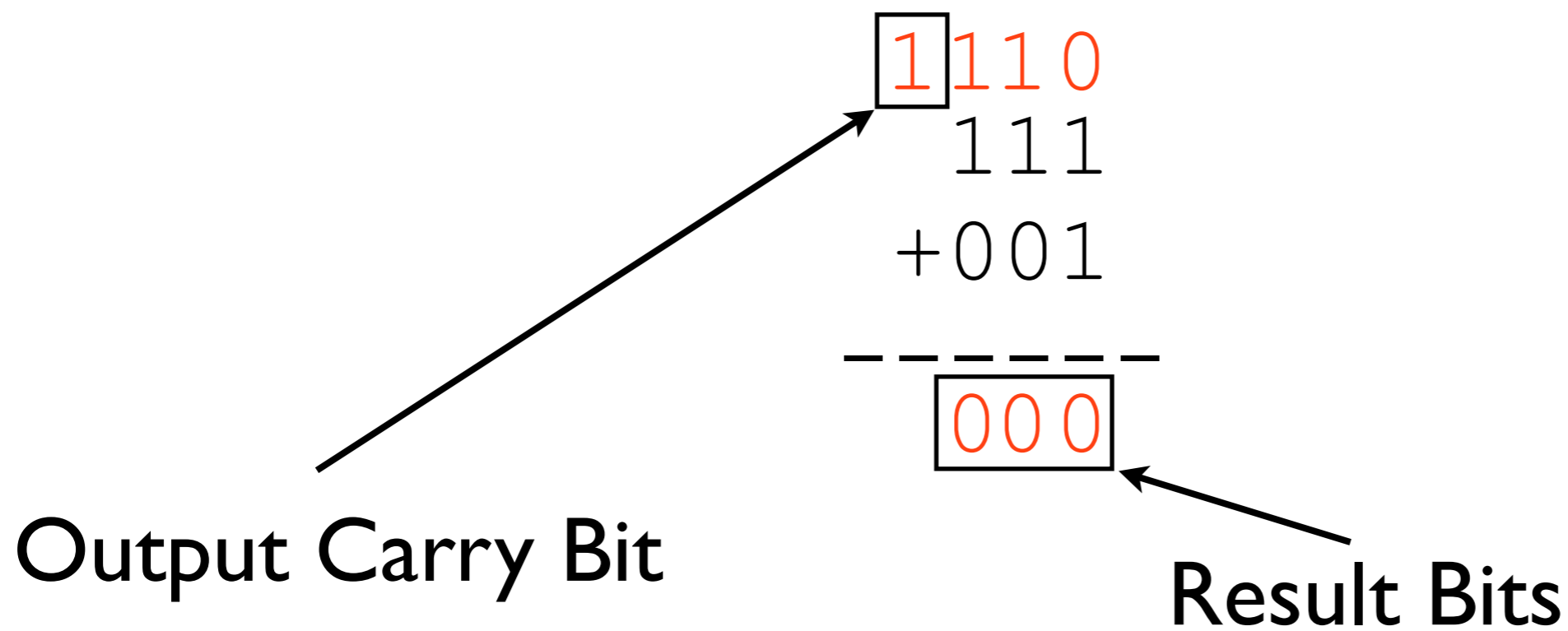
Another Example

$$\begin{array}{r} 10 \\ 111 \\ +001 \\ \hline 0 \end{array}$$

Another Example

$$\begin{array}{r} 110 \\ 111 \\ +001 \\ \hline 00 \end{array}$$

Another Example



-Now we have an output carry bit of 1. What does this mean?

Output Carry Bit Significance

- For unsigned numbers, it indicates if the result did not fit all the way into the number of bits allotted
- May be an error condition for software

Signed Addition

- Question: what is the result of the following operation?

$$\begin{array}{r} 011 \\ +011 \\ \hline \end{array}$$

?

Signed Addition

- Question: what is the result of the following operation?

$$\begin{array}{r} 011 \\ +011 \\ \hline 0110 \end{array}$$

- If these are treated as signed numbers in two's complement, then we need a leading 0 to indicate that this is a positive number
- Truncated to three bits, the result is a negative number!

Overflow

- In this situation, *overflow* occurred: this means that both the operands had the same sign, and the result's sign differed

$$\begin{array}{r} 011 \\ +011 \\ \hline 110 \end{array}$$

- Possibly a software error

Overflow vs. Carry

- These are **different ideas**
 - Carry is relevant to **unsigned** values
 - Overflow is relevant to **signed** values

```
  111
+001
----
  000
```

No Overflow;
Carry

```
  011
+011
----
  110
```

Overflow;
No Carry

```
  111
+100
----
  011
```

Overflow;
Carry

```
  001
+001
----
  010
```

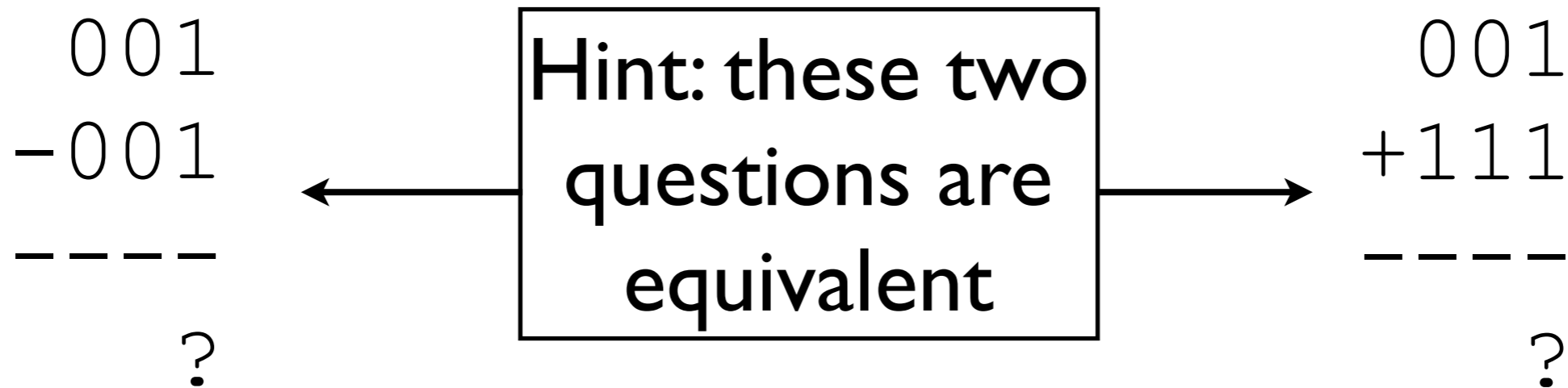
No Overflow;
No Carry

-As to when is it a problem, this all depends on exactly what it is you're doing

Subtraction

Subtraction

- Have been saying to invert bits and add one to second operand
- Could do it this way in hardware, but there is a trick



Subtraction Trick

- Assume we can cheaply invert bits, but we want to avoid adding twice (once to add 1 and once to add the other result)
- How can we do this easily?

Subtraction Trick

- Assume we can cheaply invert bits, but we want to avoid adding twice (once to add 1 and once to add the other result)
- How can we do this easily?
 - Set the initial carry to 1 instead of 0

Subtraction Example

```
  0101  
- 0011  
-----
```

Subtraction Example

0101
-0011

Invert 0011
—————→

Subtraction Example

$$\begin{array}{r} 0101 \\ -0011 \\ \hline \end{array} \quad \begin{array}{l} \text{Invert } 0011 \\ \hline \end{array} \quad \begin{array}{r} 1100 \end{array}$$

Subtraction Example

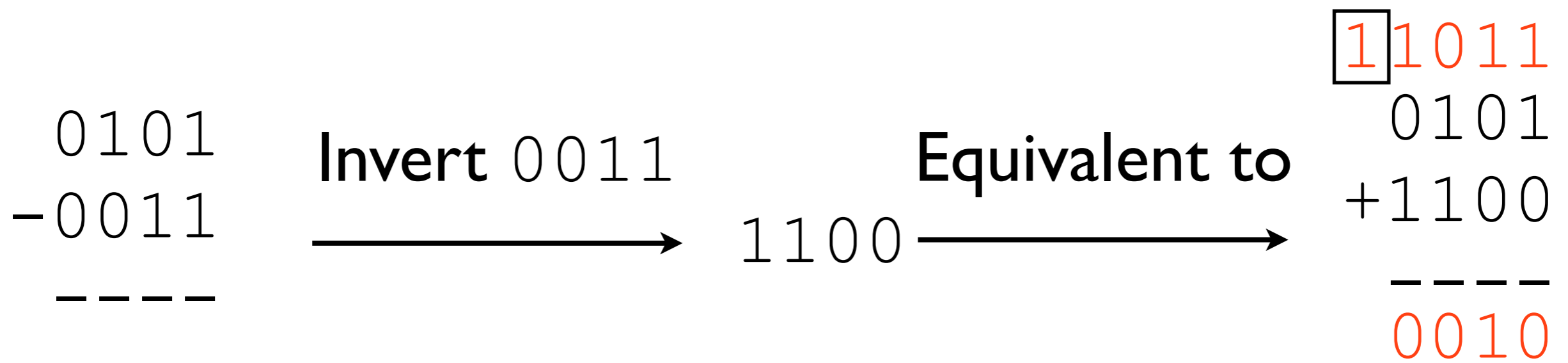
$$\begin{array}{r} 0101 \\ -0011 \\ \hline \end{array} \quad \begin{array}{l} \text{Invert } 0011 \\ \longrightarrow \end{array} \quad 1100 \quad \begin{array}{l} \text{Equivalent to} \\ \longrightarrow \end{array}$$

Subtraction Example

$$\begin{array}{r} 0101 \\ -0011 \\ \hline \end{array} \xrightarrow{\text{Invert } 0011} 1100 \xrightarrow{\text{Equivalent to}} \begin{array}{r} 0101 \\ +1100 \\ \hline \end{array} \begin{array}{c} 1 \\ 1 \\ 1 \\ 1 \end{array}$$

-An initial carry-in of 1 is equivalent to adding 1 and then adding the other operand

Subtraction Example



-An initial carry-in of 1 is equivalent to adding 1 and then adding the other operand