COMP 410 Lecture 2

Kyle Dewey

SAT and Semantic Tableau

SAT Background

SAT

- Short for the Boolean satisfiability problem
- Given a Boolean formula with variables, is there an assignment of true/false to the variables which makes the formula true?

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Yes: x is true, z is true

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Yes: x is true, z is true

$(x \land \neg x)$

No

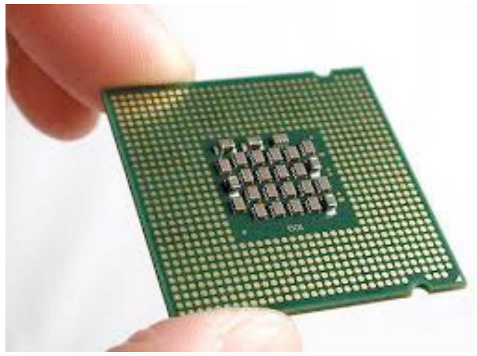
Relevance

Widespread usage in hardware and software verification

-Verification as in _proving_ the system does what we intend
-Much stronger guarantees than testing
-Testing can prove the existence of a bug (a failed test), whereas verification proves the absence of bugs (relative to the theorems proven)

Relevance

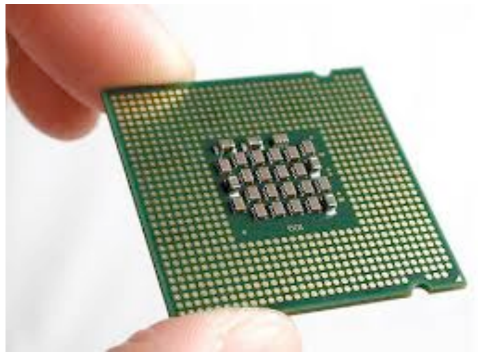
Widespread usage in hardware and software verification



-Circuits can be represented as Boolean formulas
-Can basically phrase proofs as Circuit $\wedge$ BadThing. If unsatisfiable, then BadThing cannot occur. If satisfiable, then the solution gives the circumstance under which BadThing occurs.
-Many details omitted (entire careers are based on this stuff)

Relevance

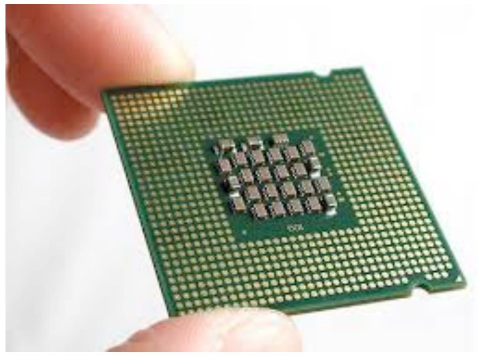
Widespread usage in hardware and software verification



-(Likely) used by AirBus to verify that flight control software does the right thing
-Lots of proprietary details so it’s not 100% clear how this verification works, but SAT is still relevant to the problem

Relevance

Widespread usage in hardware and software verification



-Nasa uses software verification for a variety of tasks; SAT is relevant, though other techniques are used, too

Relevance to Logic Programming

- Methods for solving SAT can be used to execute logic programs
- Logic programming can be phrased as SAT with some additional stuff

Semantic Tableau

- One method for solving SAT instances
- Basic idea: iterate over the formula
  - Maintain subformulas that must be true
  - Learn which variables must be true/false
  - Stop at conflicts (unsatisfiable), or when no subformulas remain (have solution)

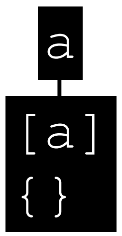
–There are many methods to this

Positive Literals

a

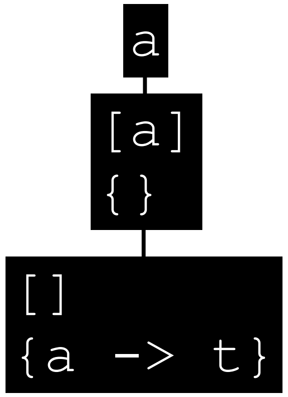
–As in, the input formula is simply “a”

Positive Literals



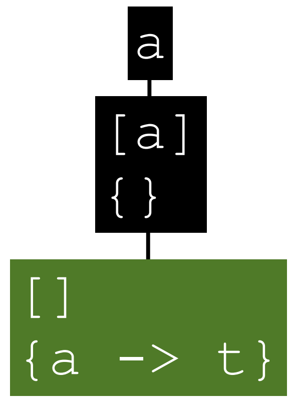
-One subformula must be true: a
-Initially, we don't know what any variables must map to

Positive Literals



-For formula “a” to be true, it must be the case that a is true

Positive Literals



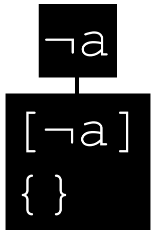
-No subformulas remain, so we are done. The satisfying solution is that a must be true.

Negative Literals

¬a

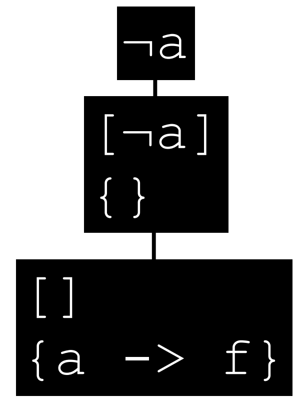
–As in, the input formula is simply "¬a"

Negative Literals



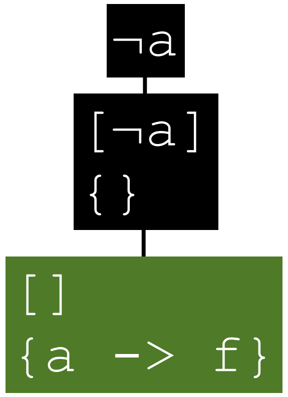
-One subformula must be true: ¬a
-Initially, we don't know what any variables must map to

Negative Literals



-For subformula “¬a” to be true, it must be the case that a is false

Negative Literals

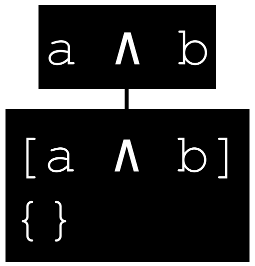


-No subformulas remain, so we are done. The satisfying solution is that “a” must be false.

Logical And

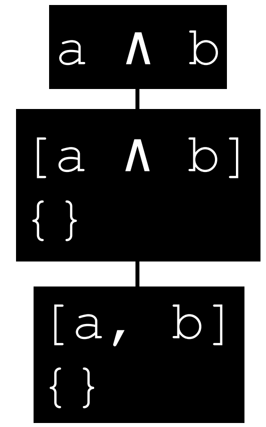
$a \land b$

Logical And



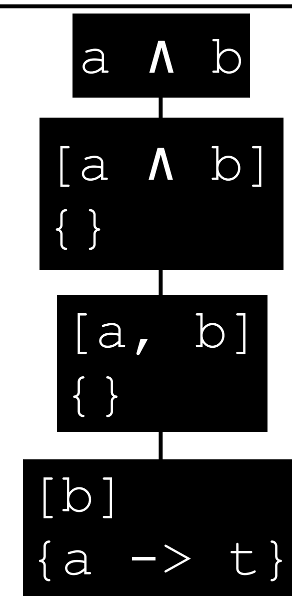
–Initially, one subformula must be true: a ∧ b
–Initially, we don’t know what any variable must map to

Logical And



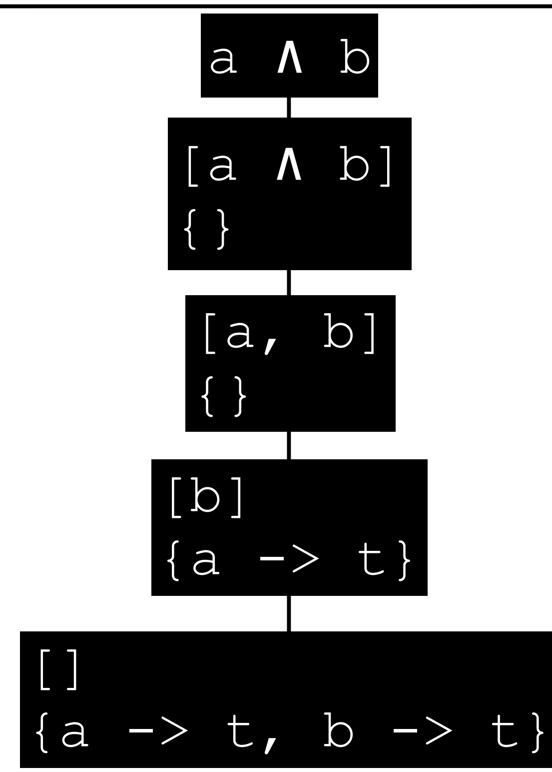
-For a ∧ b to be true, subformulas a and b must both be true

Logical And



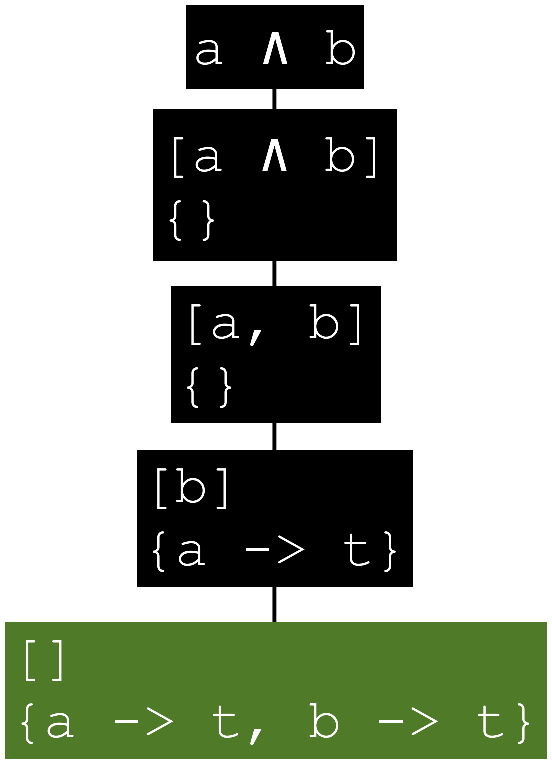
-From the positive literal case, for formula a to be true, variable a must be true

Logical And



-From the positive literal case, for formula b to be true, variable b must be true

Logical And



-No subformulas remain, so we are done with the solution that both a and b must be true

Logical And

$a \wedge \neg a$

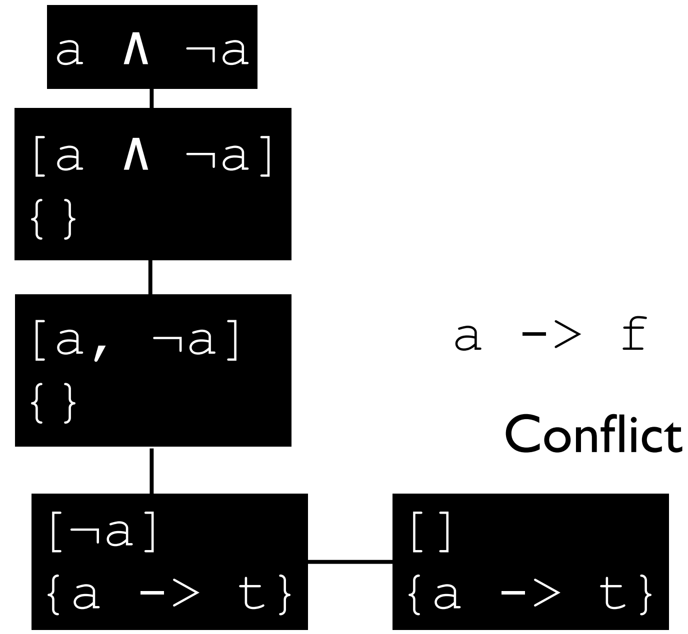
–Alternative example, showing a conflict

Logical And

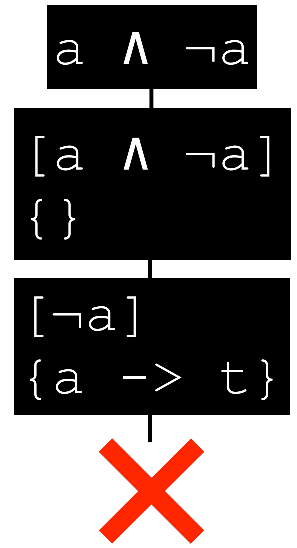
a ∧ ¬a

[a ∧ ¬a]
{}

Logical And



Logical And



-Now we have a problem: for formula ¬a to be true, it must be the case that variable a is false
-We've already recorded that variable a must be true, which is the opposite of what we expect.
-As such, we have a conflict - this formula is unsatisfiable

Exercise: First Side of SAT Sheet

Logical Or

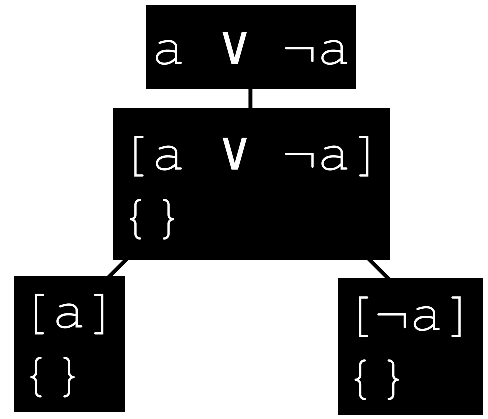
$a \lor \neg a$

Logical Or

```
a ∨ ¬a
```

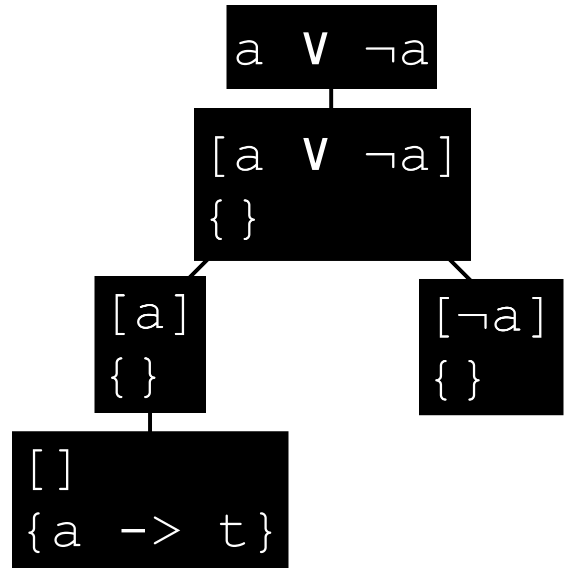
```
[a ∨ ¬a]
{}
```

Logical Or



-World splits on or: in one world we pick the left subformula, and in another we pick the right

Logical Or



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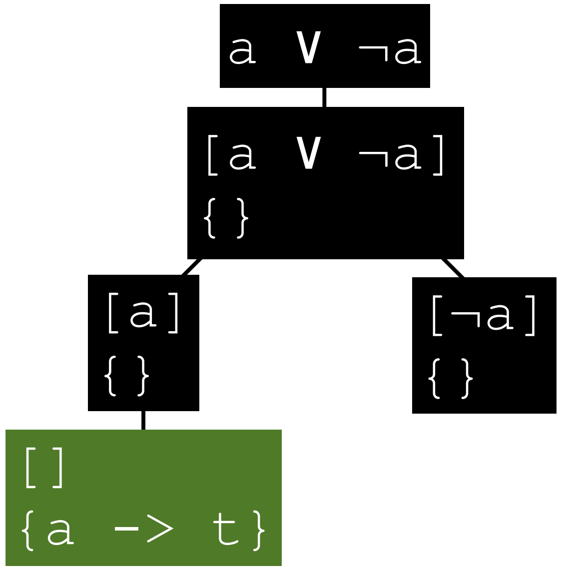
Logical Or



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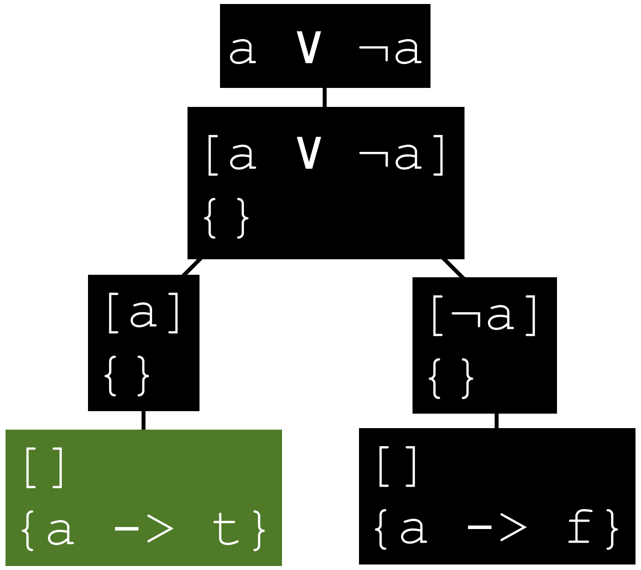
Logical Or



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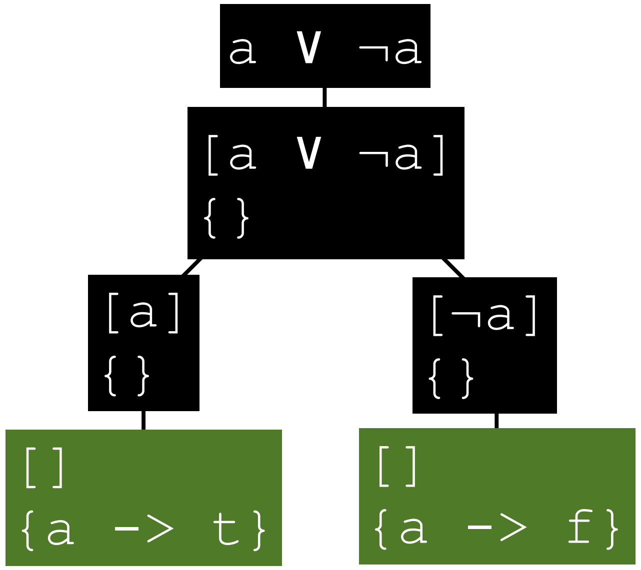
Logical Or



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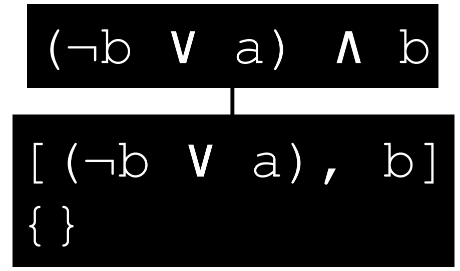
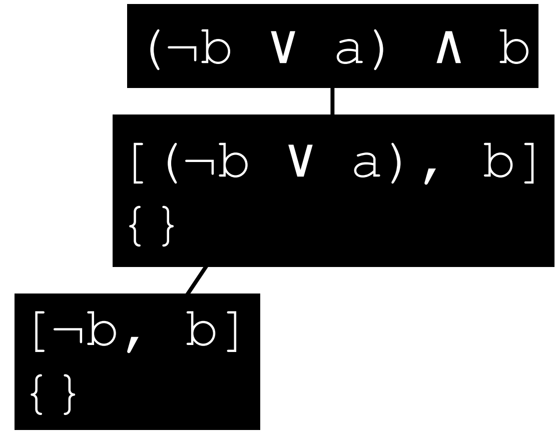
Examples

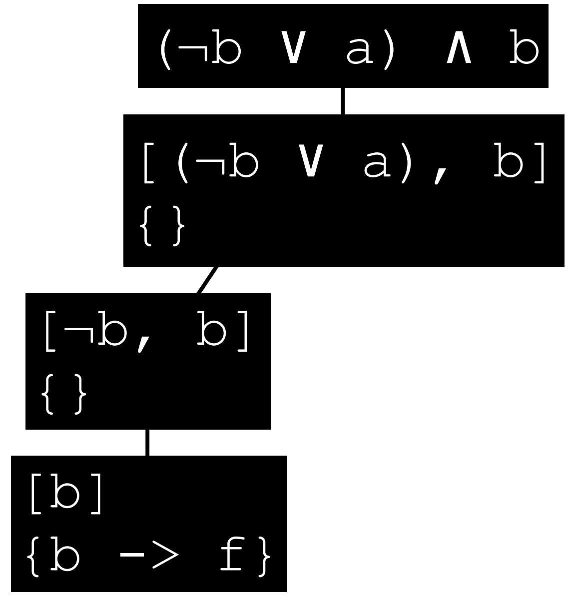
Example 1:

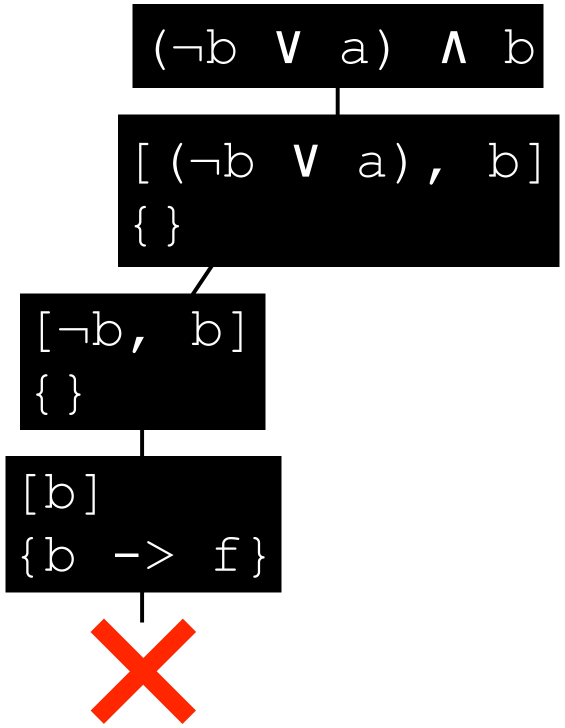
$(\neg b \lor a) \land b$

$(\neg b \lor a) \land b$


(¬b ∨ a) ∧ b
[(¬b ∨ a), b]
{}


(¬b ∨ a) ∧ b
[(¬b ∨ a), b]
{}
[¬b, b]
{}


(¬b ∨ a) ∧ b
[(¬b ∨ a), b]
{}
[¬b, b]
{}
[b]
{b -> f}



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[¬b, b]
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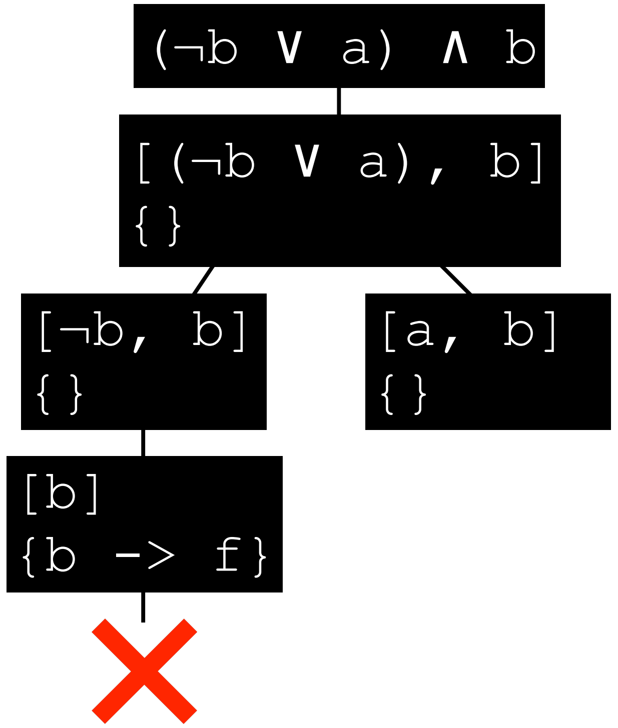
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[a, b]
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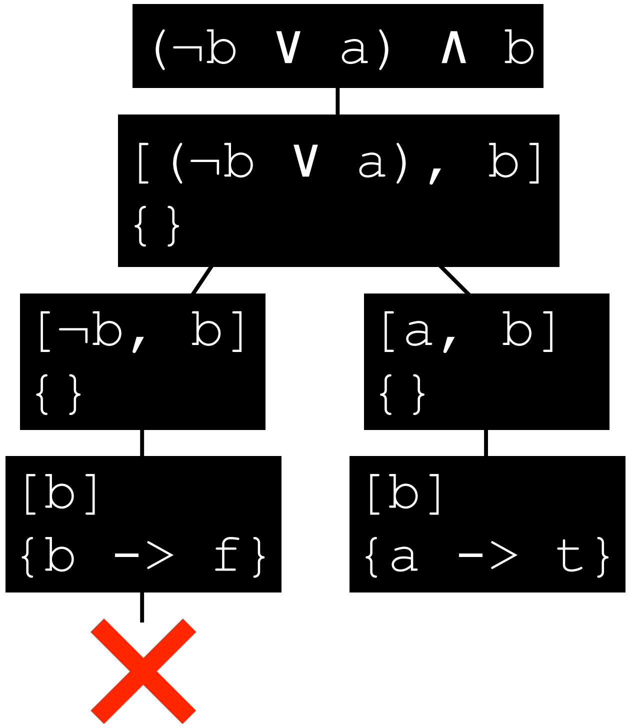
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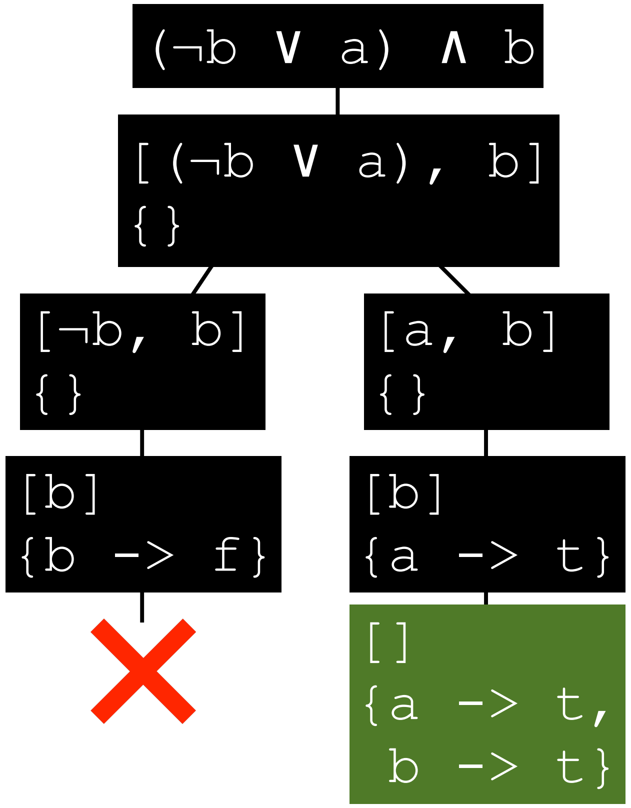
[(¬b ∨ a), b]
{}
[¬b, b]
{}
[a, b]
{}
[b]
{b -> f}
[b]
{a -> t}


(¬b ∨ a) ∧ b
[(¬b ∨ a), b]
{}
[¬b, b]
{}
[a, b]
{}
[b]
{b -> f}
[b]
{a -> t}
[]
{a -> t,
b -> t}

Example 2:

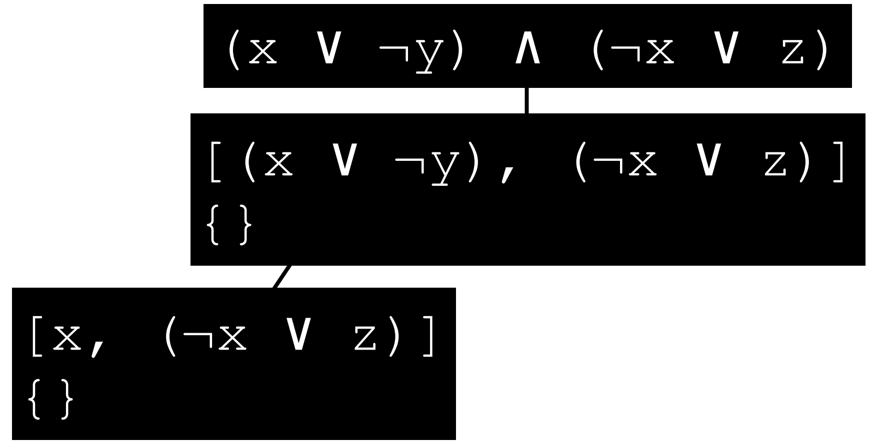
$(x \lor \neg y) \land (\neg x \lor z)$

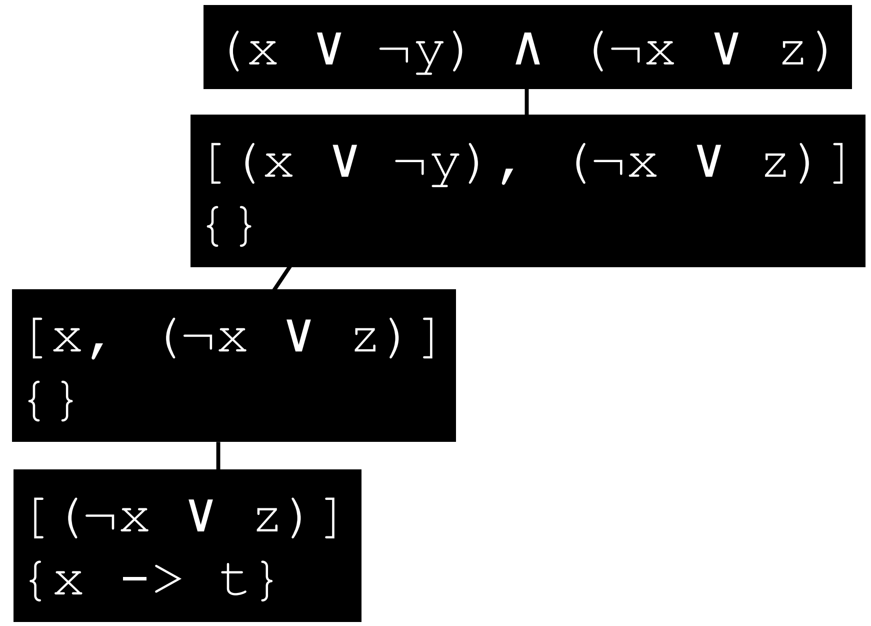
$(x \lor \neg y) \land (\neg x \lor z)$

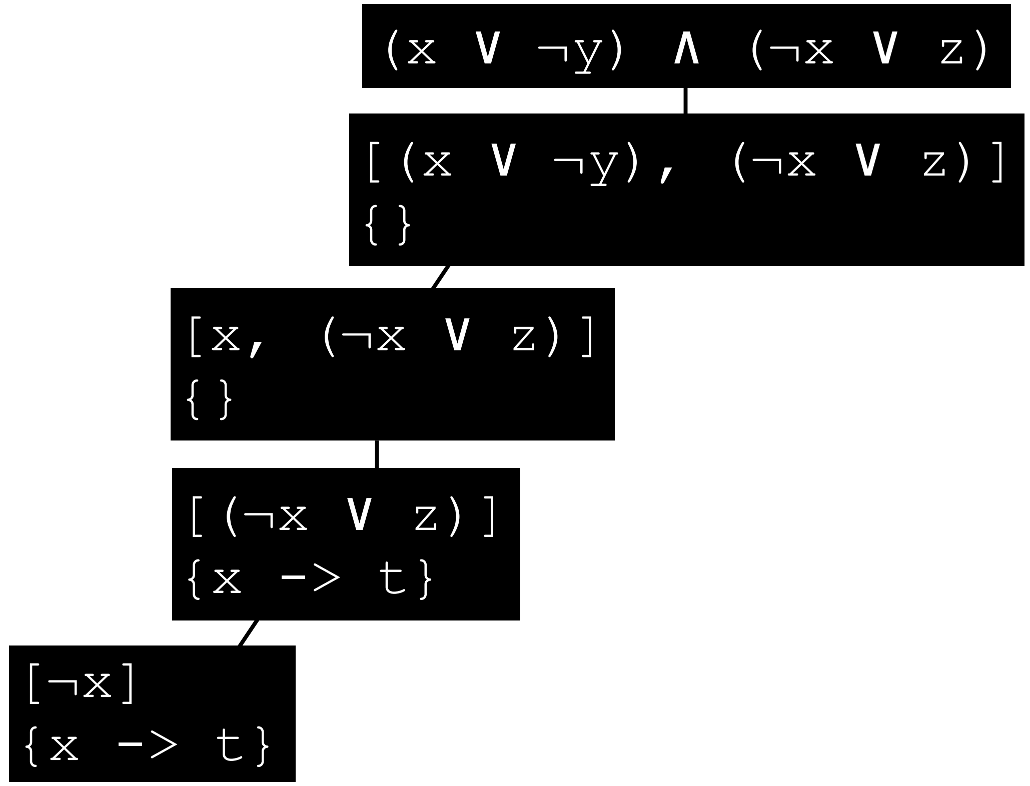
$(x \lor \neg y) \land (\neg x \lor z)$

$[(x \lor \neg y), (\neg x \lor z)]$

$\{\}$


(x ∨ ¬y) ∧ (¬x ∨ z)
[(x ∨ ¬y), (¬x ∨ z)]
{}
[x, (¬x ∨ z)]
{}


(x ∨ ¬y) ∧ (¬x ∨ z)
[(x ∨ ¬y), (¬x ∨ z)]
{}
[x, (¬x ∨ z)]
{}
[(¬x ∨ z)]
{x -> t}


(x ∨ ¬y) ∧ (¬x ∨ z)
[(x ∨ ¬y), (¬x ∨ z)]
{}
[x, (¬x ∨ z)]
{}
[(¬x ∨ z)]
{x -> t}
[¬x]
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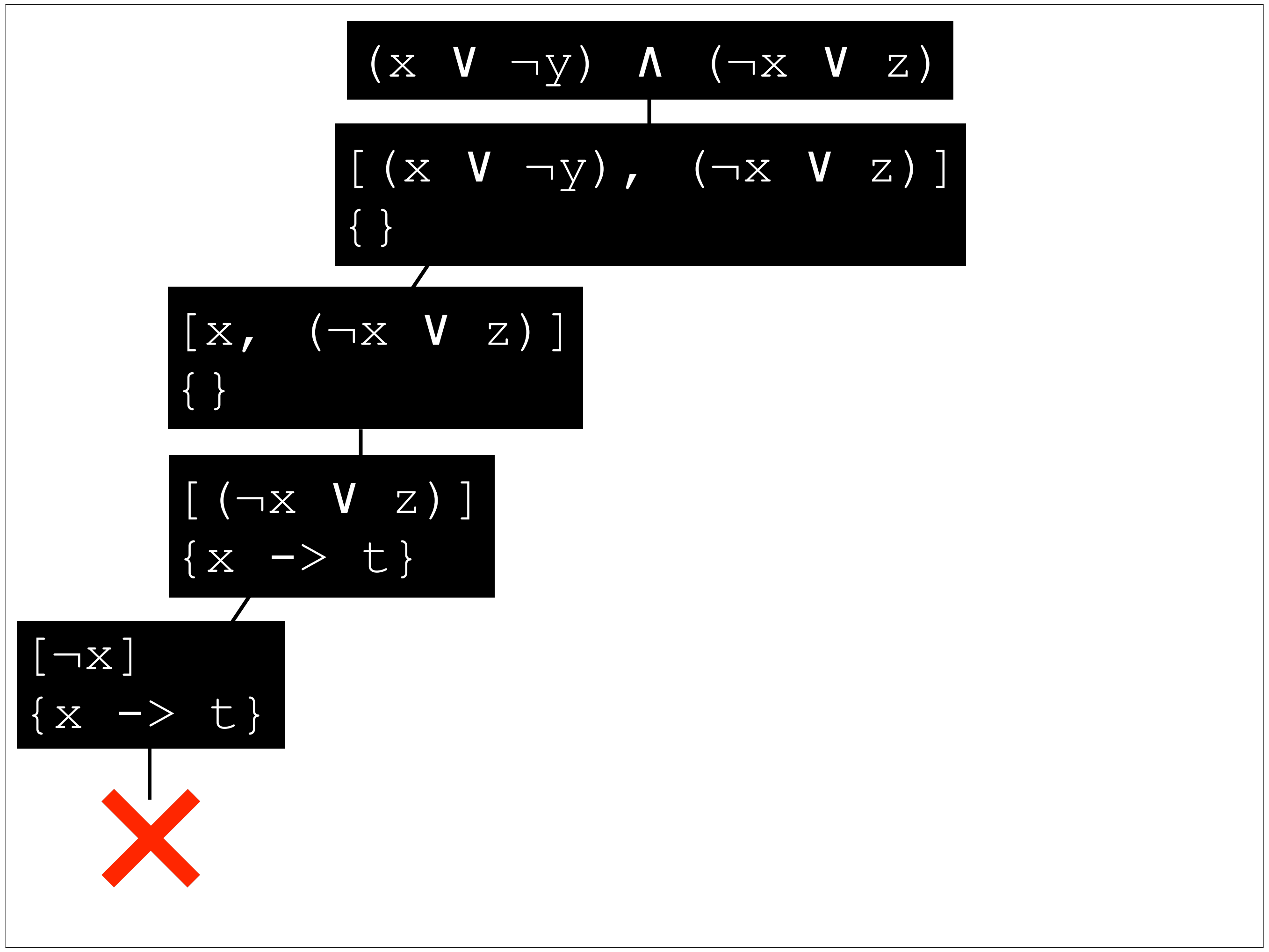
(x ∨ ¬y) ∧ (¬x ∨ z)
[(x ∨ ¬y), (¬x ∨ z)]
{}
[x, (¬x ∨ z)]
{}
[(¬x ∨ z)]
{x -> t}
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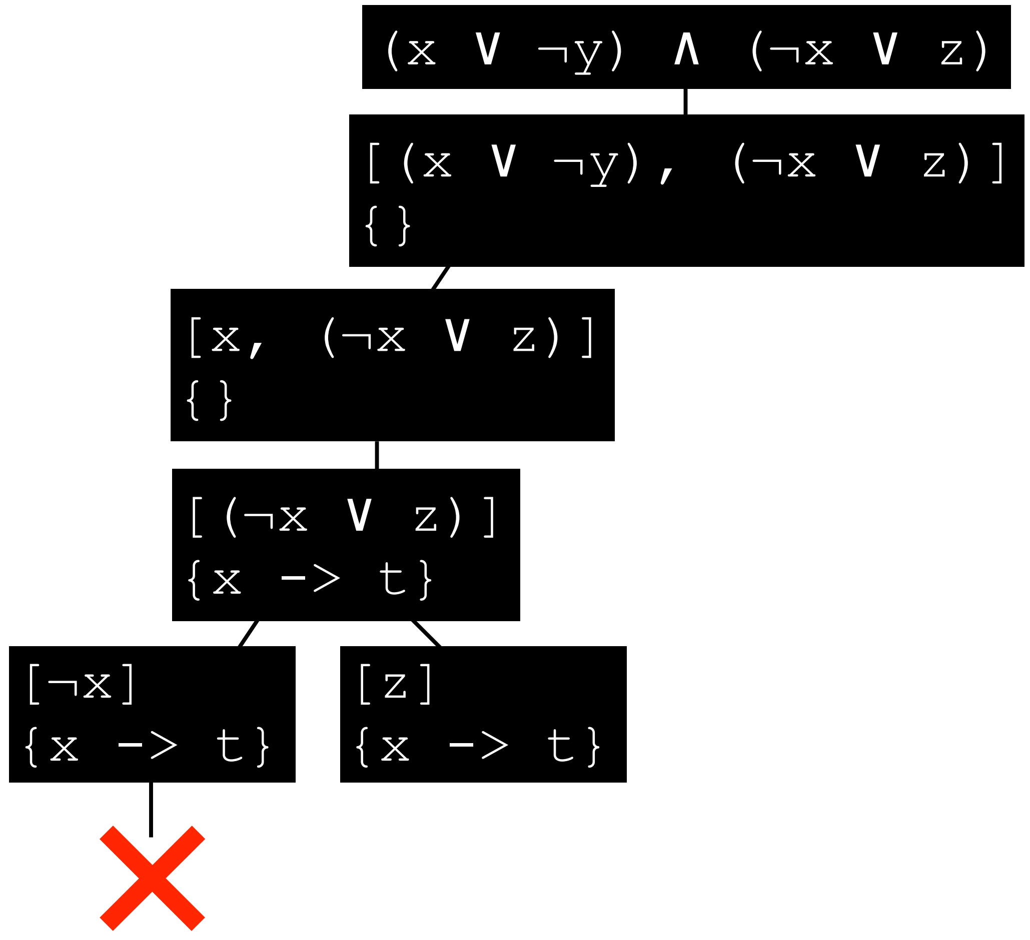
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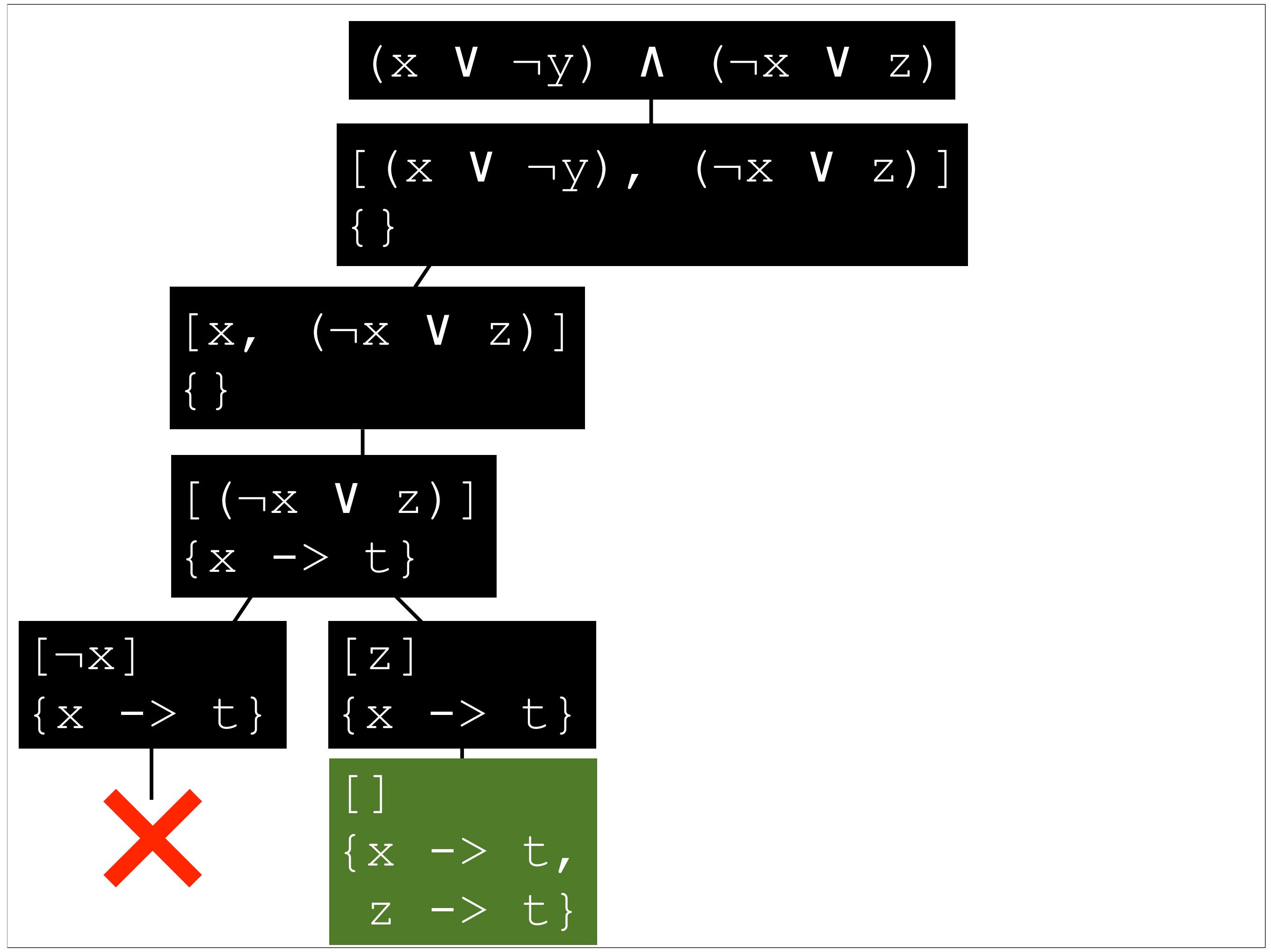
[(x ∨ ¬y), (¬x ∨ z)]
{}
[x, (¬x ∨ z)]
{}
[(¬x ∨ z)]
{x -> t}
[¬x]
{x -> t}
[z]
{x -> t}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{x -> t}
[¬x]
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[z]
{x -> t}
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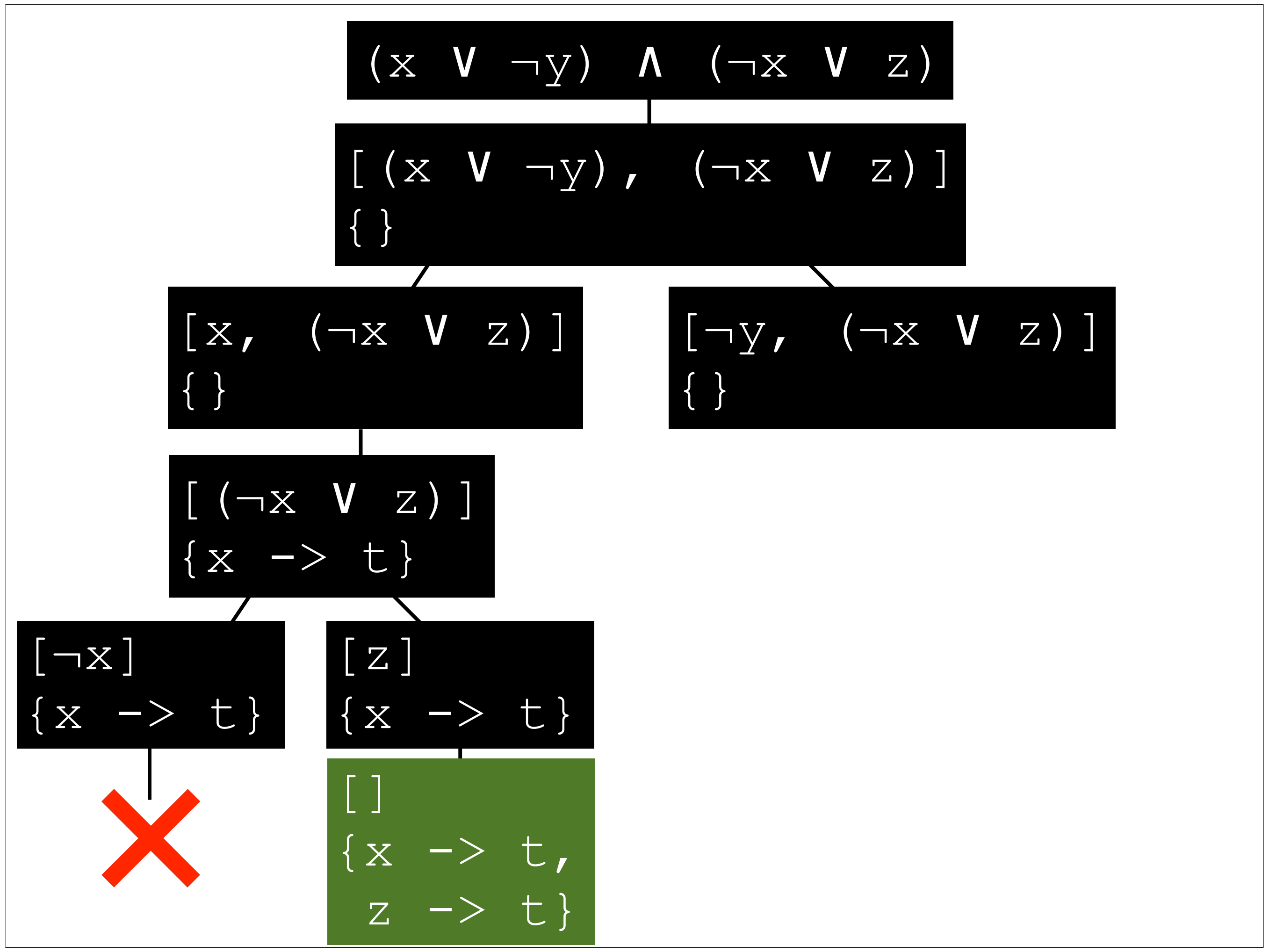
(x ∨ ¬y) ∧ (¬x ∨ z)
[(x ∨ ¬y), (¬x ∨ z)]
{}
[x, (¬x ∨ z)]
{}
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[(¬x ∨ z)]
{x -> t}
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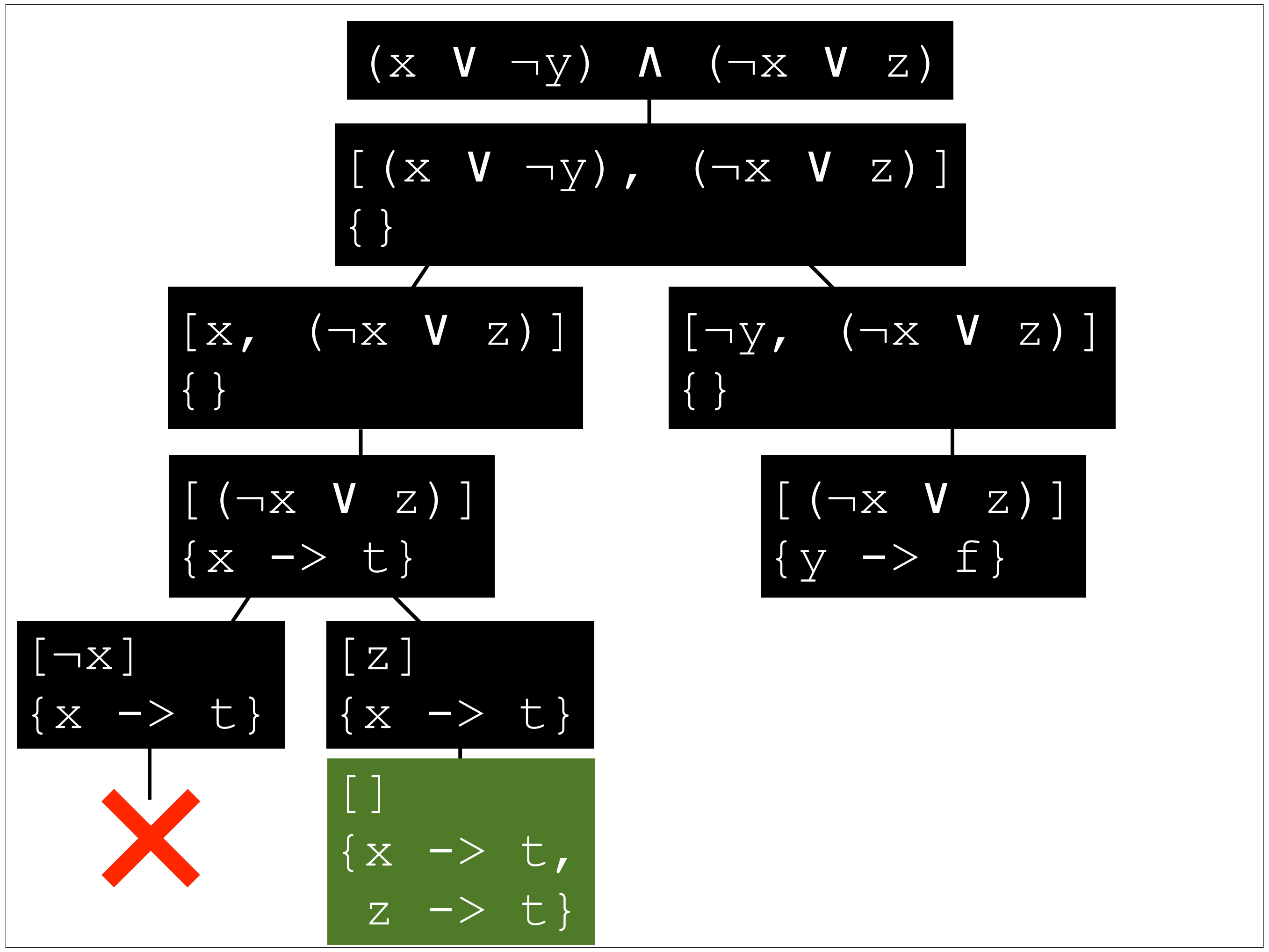
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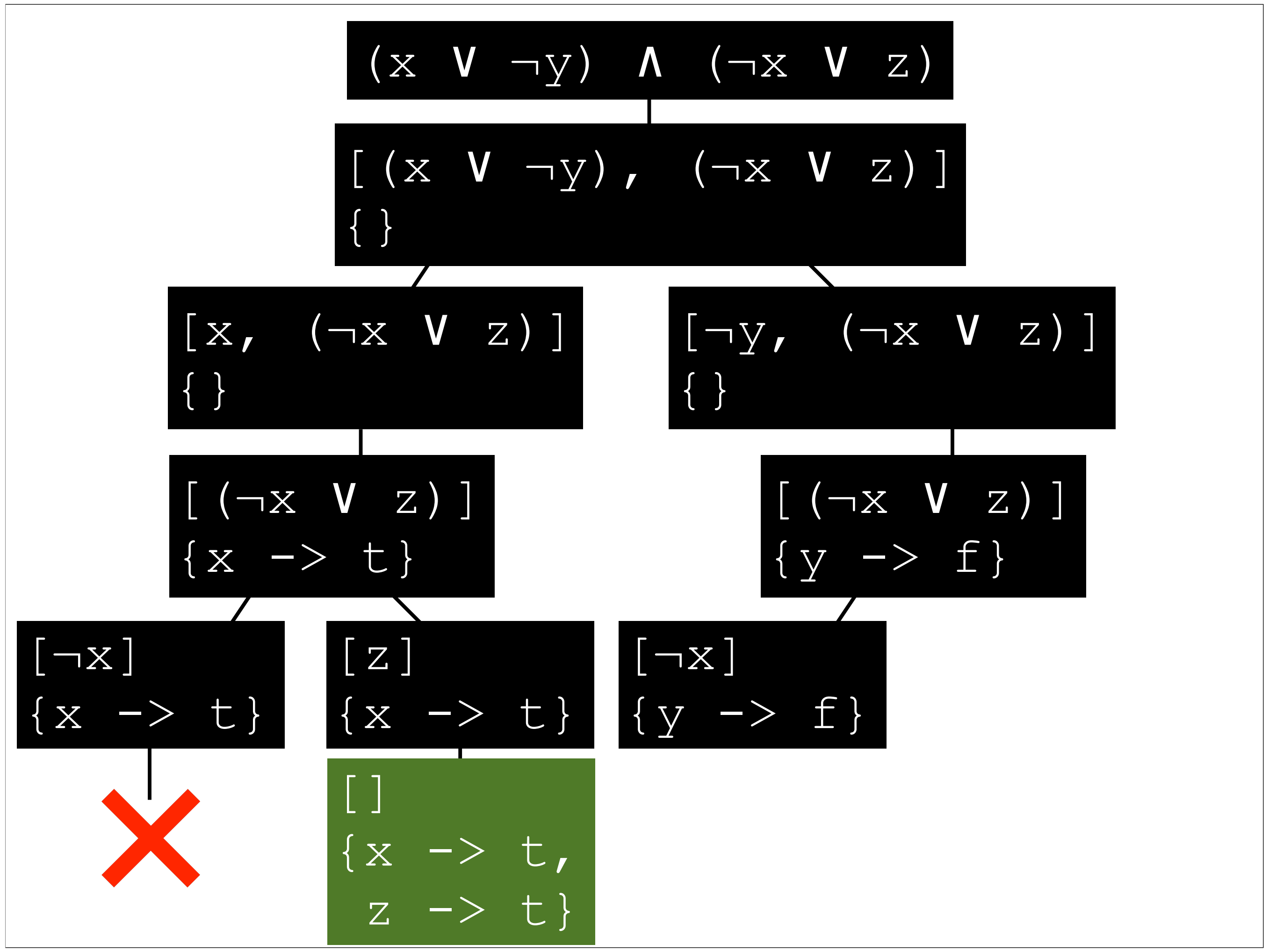
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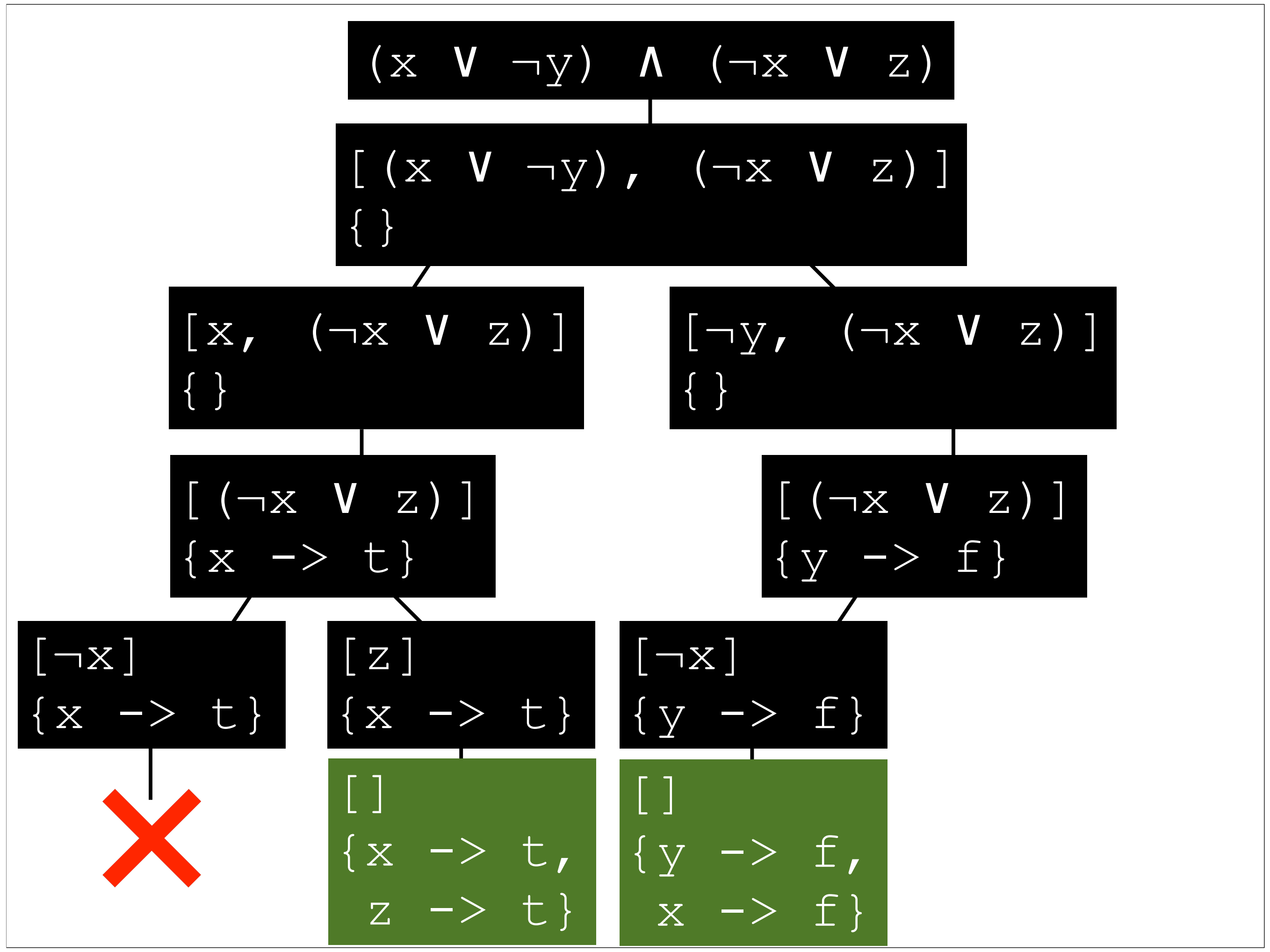
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[¬y, (¬x ∨ z)]
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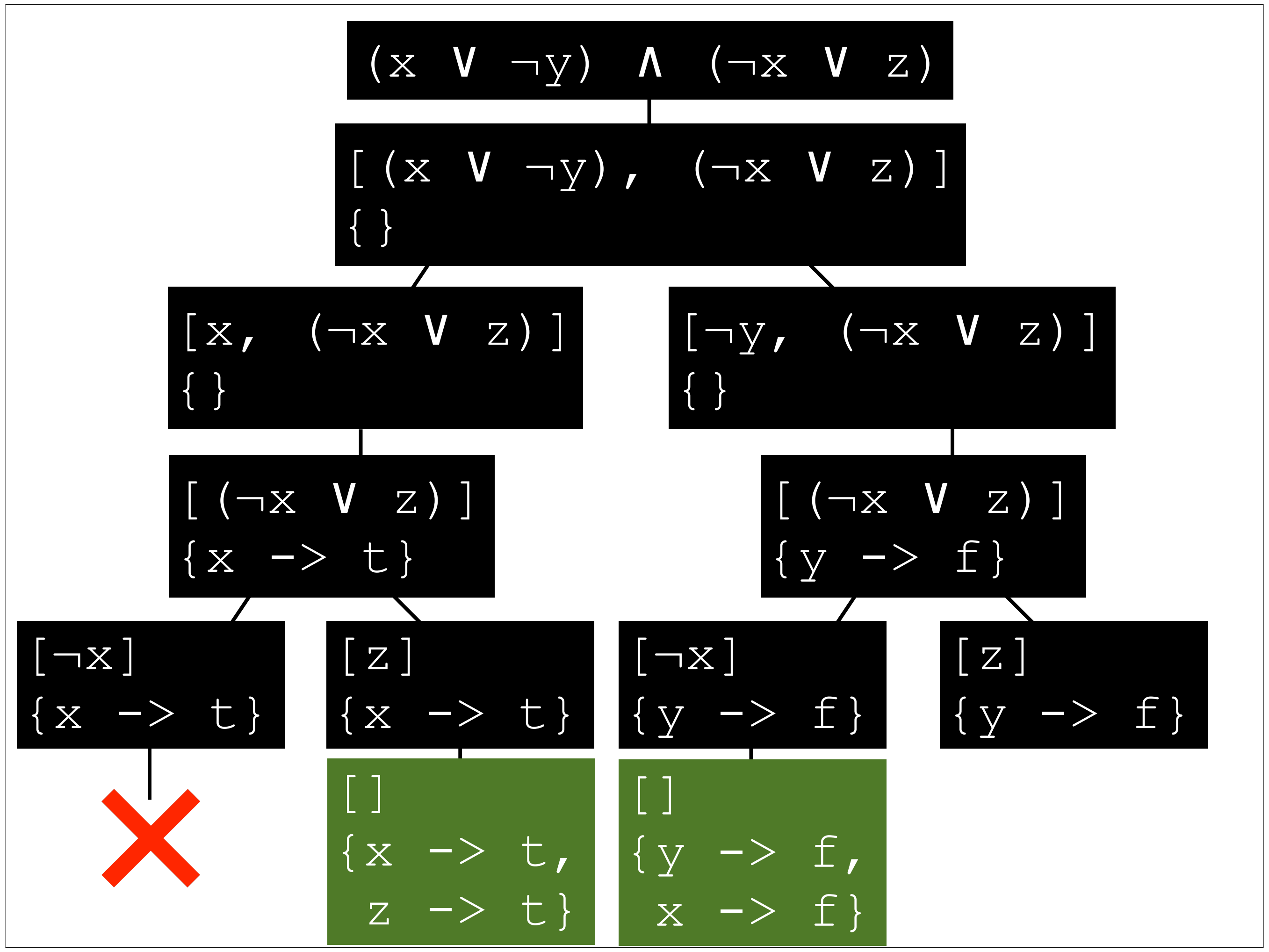
[(¬x ∨ z)]
{y -> f}
[¬x]
{x -> t}
[z]
{x -> t}
[¬x]
{y -> f}
[]
{x -> t,
z -> t}
[]
{y -> f,
x -> f}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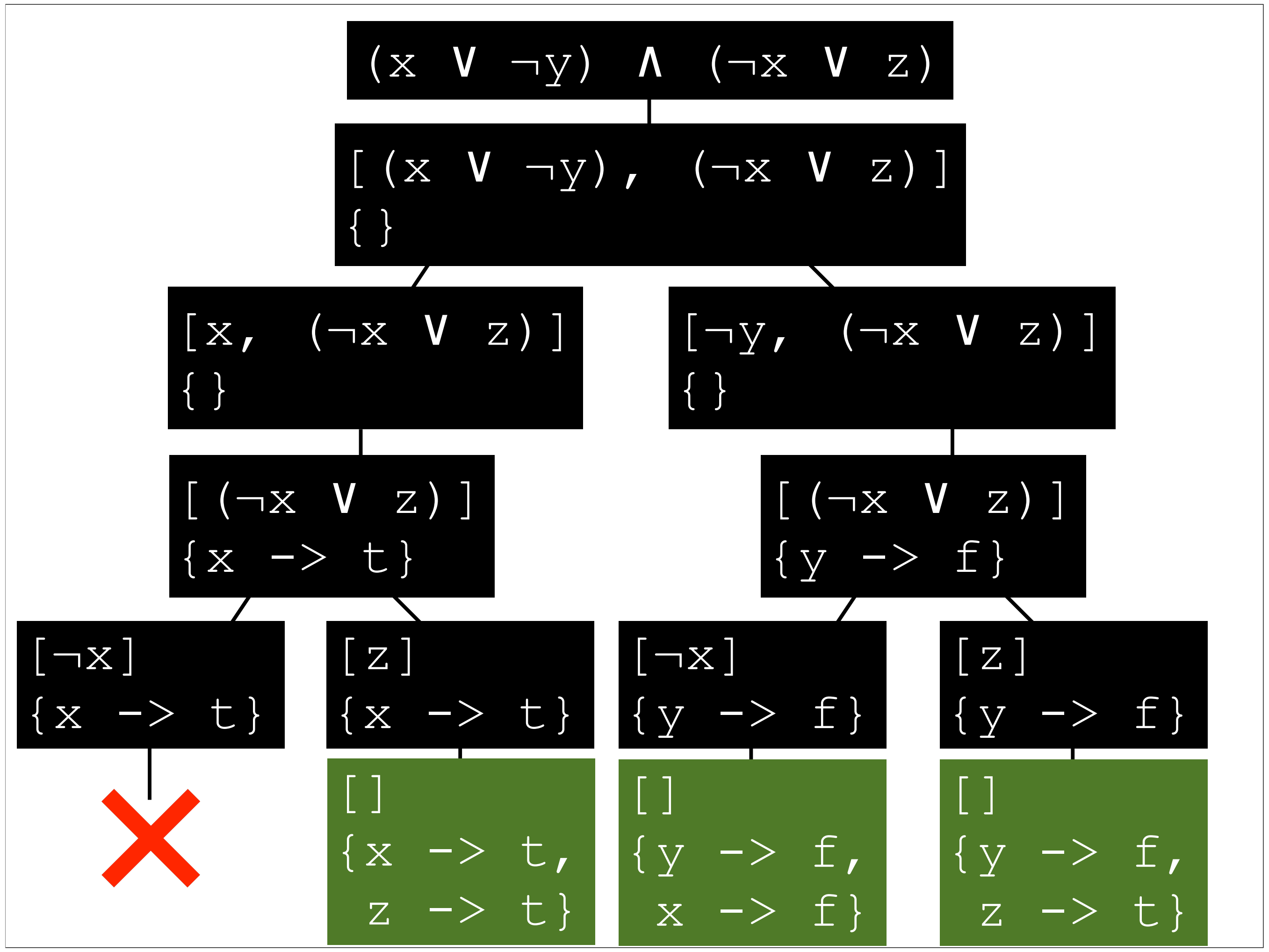
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{y -> f}
[]
{x -> t,
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Exercise: Second Side of SAT Sheet