

CS24 Week 6 Lecture 1

Kyle Dewey

Overview

- Complexity and complexity analysis

Complexity

Complexity

- Up until this point, we have used terms like “efficiency”, “expensive”, and “cheap”

```
int foo(int* array) {  
    return array[2] * array[3];  
}
```

```
int bar(int* array) {  
    int x;  
    for(x = 0; x < MAX_SIZE; x++) {  
        if (array[x] == 7) return x;  
    }  
    return -1;  
}
```

Complexity

- Up until this point, we have used terms like “efficiency”, “expensive”, and “cheap”

```
int foo(int* array) {  
    return array[2] * array[3];    cheap(?)  
}
```

```
int bar(int* array) { expensive(?)  
    int x;  
    for(x = 0; x < MAX_SIZE; x++) {  
        if (array[x] == 7) return x;  
    }  
    return -1;  
}
```

Ambiguous Terms

- Under what circumstances is this cheap?
- When is it expensive?

```
int bar(int* array) {
    int x;
    for(x = 0; x < MAX_SIZE; x++) {
        if (array[x] == 7) return x;
    }
    return -1;
}
```

“Expensive”, “Cheap”, “Efficient”

- What is good about these terms?
- What is bad about these terms?

“Expensive”, “Cheap” “Efficient”

- What is good about these terms?
 - Easy to understand
- What is bad about these terms?
 - Imprecise
 - Binary in nature (either cheap or expensive)
 - Program efficiency is often dependent on **input size**

Measuring Efficiency

- How might we determine the efficiency of a program?

Measuring Efficiency

- How might we determine the efficiency of a program?
- Benchmarks tend to be too specific (new hardware? How big of inputs do we test?)
- Better approach: define a formula in terms of the input size

Big O Notation

- A formula that gives an upper bound of how expensive something is *in the worst case*, in terms of an input size N
- Which is most efficient below?

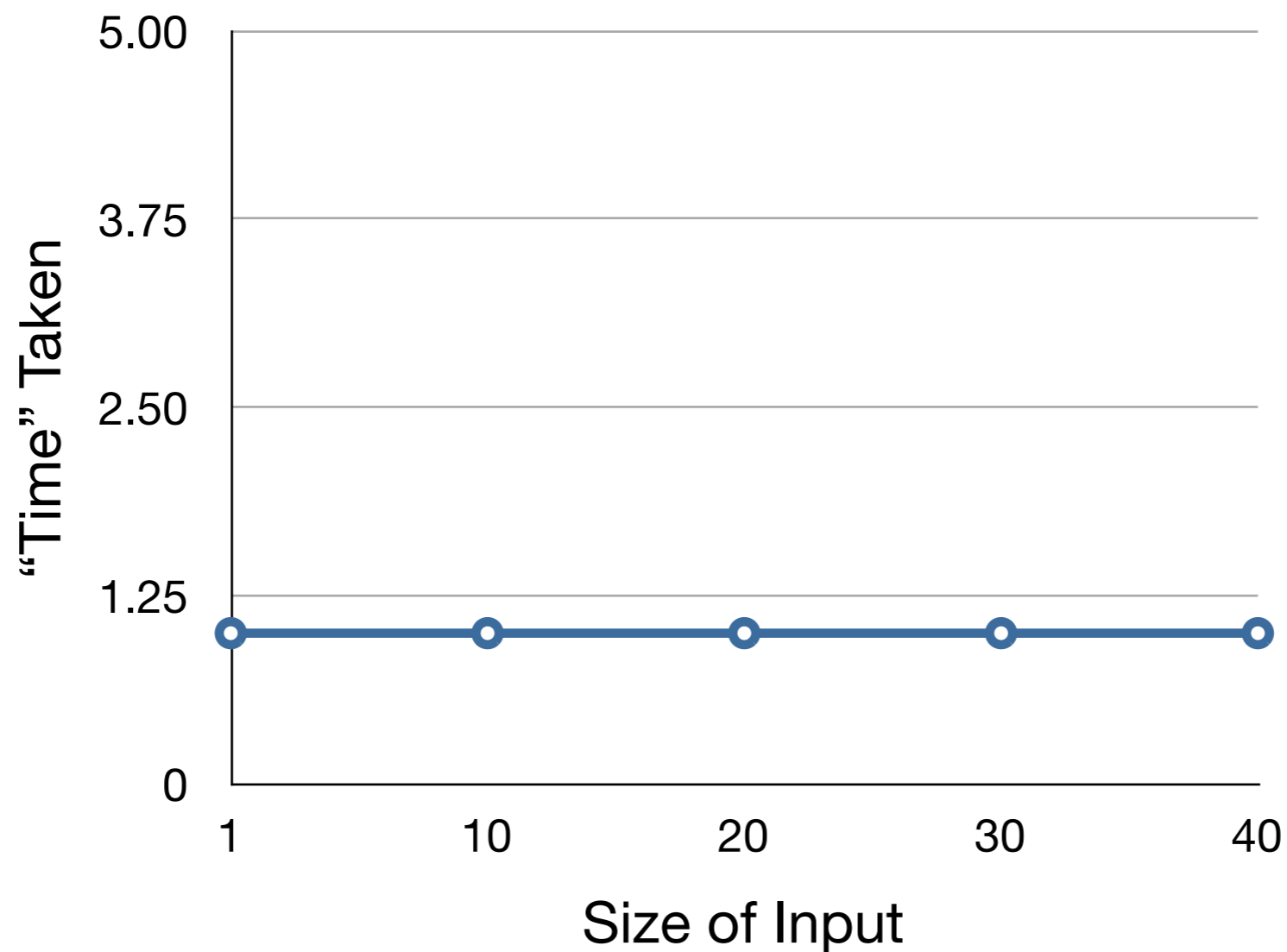
$O(1)$ // constant time

$O(n)$ // linear time

$O(n^2)$ // quadratic time

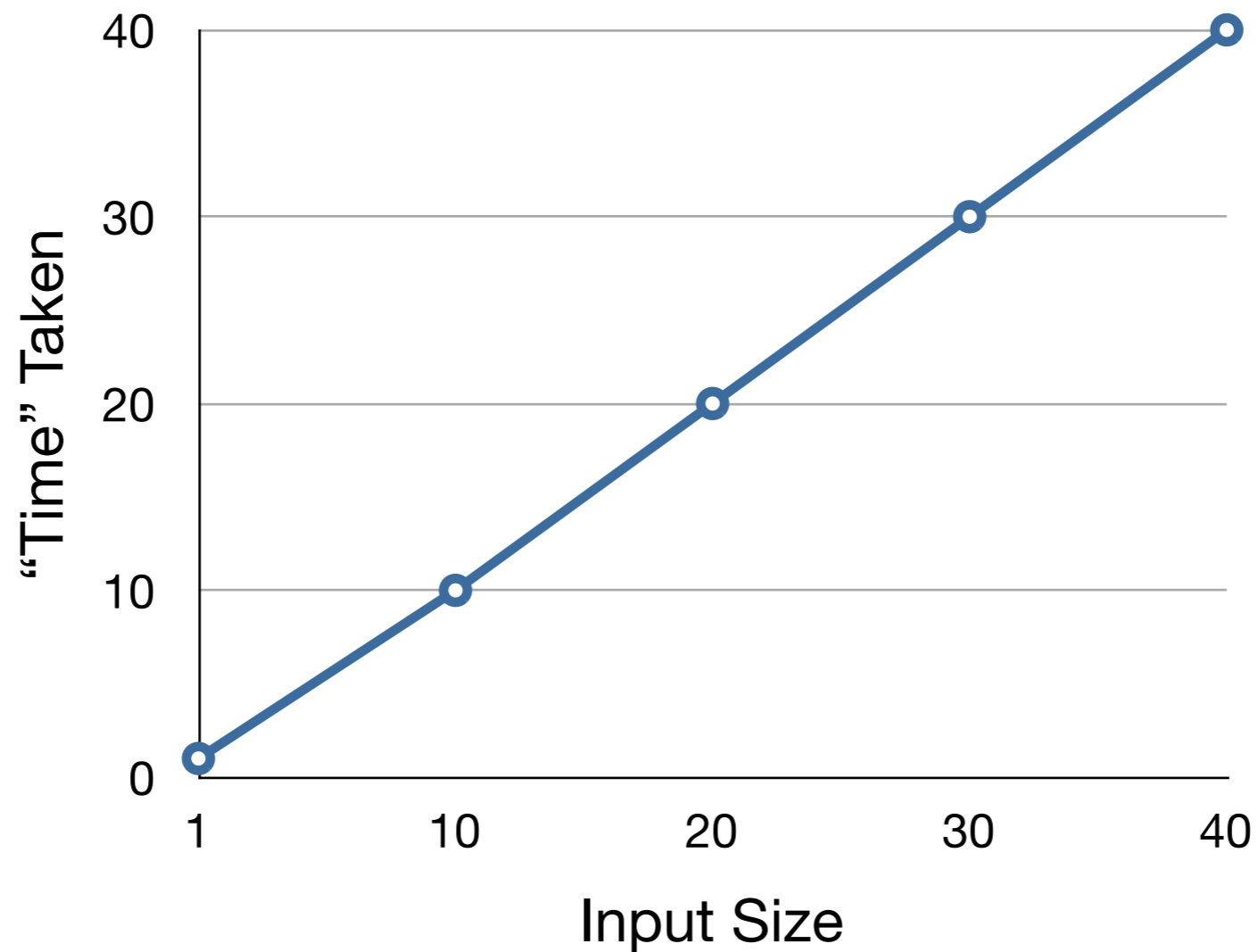
$$O(1)$$

- Regardless of the size of the input, it takes the same amount of time



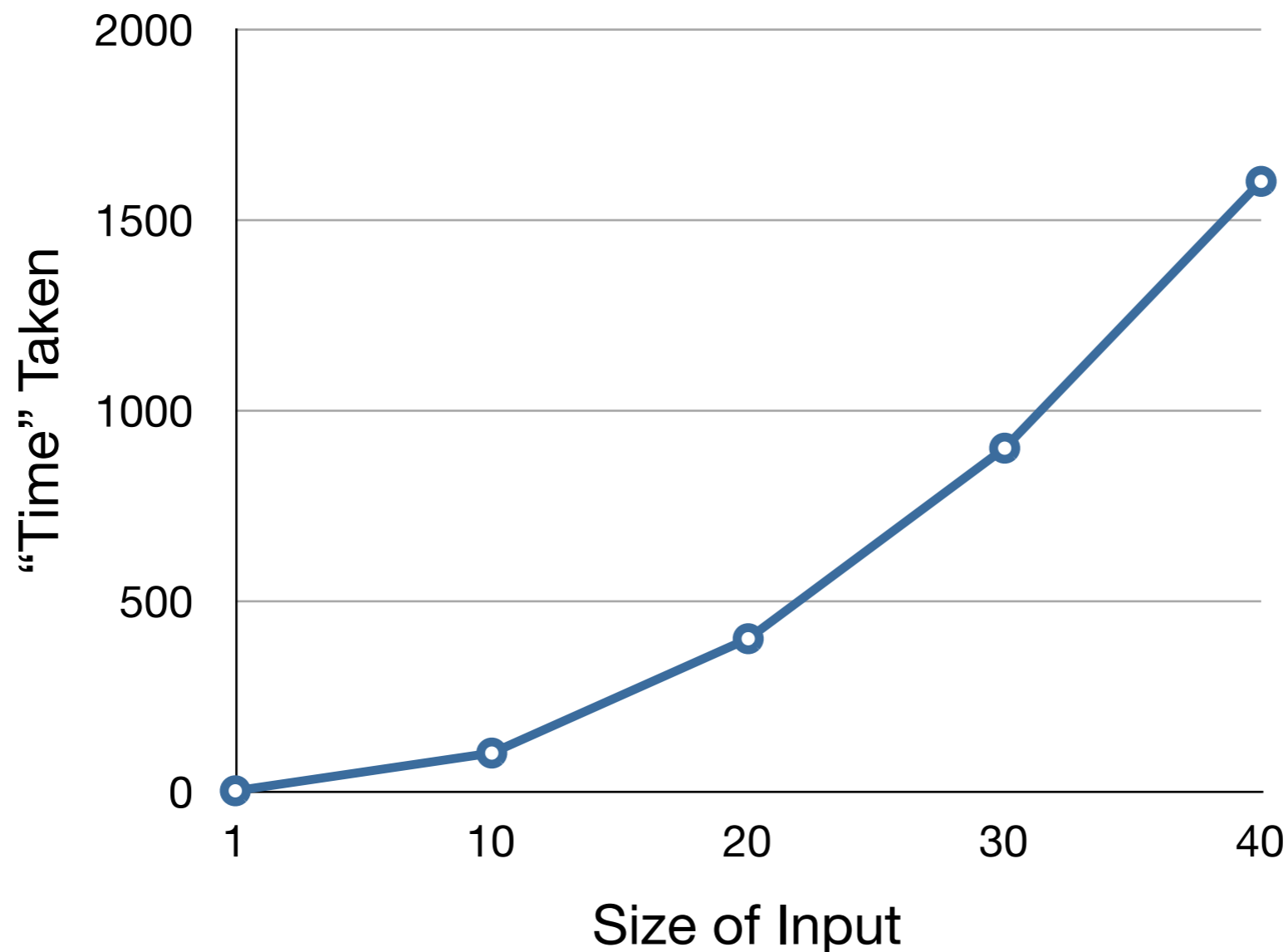
$$O(N)$$

- The amount of time taken increases linearly with the input size



$$O(n^2)$$

- The amount of time increases quadratically with input size



Determining Big O

```
int sum(int* arr, int length) {  
    int s = 0, x;  
    for (x = 0; x < length; x++) {  
        s += arr[x];  
    }  
    return s;  
}
```

Determining Big O

```
int sum(int* arr, int length) {  
    int s = 0, x;  
    for (x = 0; x < length; x++) {  
        s += arr[x];  
    }  
    return s;  
}
```

Constant time, done once. Call this C_1 .

Determining Big O

```
int sum(int* arr, int length) {  
    int s = 0, x;  
    for (x = 0; x < length; x++) {  
        s += arr[x];  
    }  
    return s;  
}
```

Constant time, done once. Call this c_2 .

Determining Big O

```
int sum(int* arr, int length) {  
    int s = 0, x;  
    for (x = 0; x < length; x++) {  
        s += arr[x];  
    }  
    return s;  
}
```

Constant time, done `length` times. Call this c_3 .

Determining Big O

```
int sum(int* arr, int length) {  
    int s = 0, x;  
    for (x = 0; x < length; x++) {  
        s += arr[x];  
    }  
    return s;  
}
```

Constant time, done `length` times. Call this c_4 .

Determining Big O

```
int sum(int* arr, int length) {  
    int s = 0, x;  
    for (x = 0; x < length; x++) {  
        s += arr[x];  
    }  
    return s;  
}
```

Constant time, done `length` times. Call this c_5 .

Determining Big O

```
int sum(int* arr, int length) {  
    int s = 0, x;  
    for (x = 0; x < length; x++) {  
        s += arr[x];  
    }  
    return s;  
}
```

Constant time, done once. Call this C_6 .

Determining Big O

- Putting it together, we get the formula:

$$c1 + c2 + (c3 * length) + (c4 * length) \\ + (c5 * length) + c6$$

Determining Big O

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$c_1 + c_2 + (c_3 * \text{length}) + (c_4 * \text{length}) \\ + (c_5 * \text{length}) + c_6$$

Determining Big O

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$1 + 1 + (1 * \text{length}) + (1 * \text{length}) + (1 * \text{length}) + 1$$

Determining Big O

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

`3 + length + length + length`

Determining Big O

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$3 + 3(\text{length})$$

Determining Big O

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

`1 + length`

Determining Big O

- With sums, we always choose the larger sum
- A variable is always larger than a constant

`1 + length`

Determining Big O

- With sums, we always choose the larger sum
- A variable is always larger than a constant

length

Determining Big O

- Observe that `length` is really N , the input size
- For this example, we are done

`length`

Determining Big O

- Observe that `length` is really N , the input size
- For this example, we are done

$O(N)$

Another Example

```
int sum2(int* arr, int length) {
    int s = 0, x, y;
    for (x = 0; x < length; x++) {
        for (y = 0; y < length; y++) {
            s += arr[x] + arr[y];
        }
    }
    return s;
}
```


Another Example

```
int sum2(int* arr, int length) {  
    int s = 0, x, y;  
    for (x = 0; x < length; x++) {  
        for (y = 0; y < length; y++) {  
            s += arr[x] + arr[y];  
        }  
    }  
    return s;  
}
```

Constant time, done once. Call this C_1 .

Another Example

```
int sum2(int* arr, int length) {  
    int s = 0, x, y;  
    for (x = 0; x < length; x++) {  
        for (y = 0; y < length; y++) {  
            s += arr[x] + arr[y];  
        }  
    }  
    return s;  
}
```

Constant time, done once. Call this C_2 .

Another Example

```
int sum2(int* arr, int length) {
    int s = 0, x, y;
    for (x = 0; x < length; x++) {
        for (y = 0; y < length; y++) {
            s += arr[x] + arr[y];
        }
    }
    return s;
}
```

Constant time, done $length$ times. Call this c_3 .

Another Example

```
int sum2(int* arr, int length) {
    int s = 0, x, y;
    for (x = 0; x < length; x++) {
        for (y = 0; y < length; y++) {
            s += arr[x] + arr[y];
        }
    }
    return s;
}
```

Constant time, done $length$ times. Call this C_4 .

Another Example

```
int sum2(int* arr, int length) {  
    int s = 0, x, y;  
    for (x = 0; x < length; x++) {  
        for (y = 0; y < length; y++) {  
            s += arr[x] + arr[y];  
        }  
    }  
    return s;  
}
```

Constant time, done length times. Call this C_5 .

Another Example

```
int sum2(int* arr, int length) {
    int s = 0, x, y;
    for (x = 0; x < length; x++) {
        for (y = 0; y < length; y++) {
            s += arr[x] + arr[y];
        }
    }
    return s;
}
```

Constant time, done $\text{length} * \text{length}$ times.
Call this C_6 .

Another Example

```
int sum2(int* arr, int length) {  
    int s = 0, x, y;  
    for (x = 0; x < length; x++) {  
        for (y = 0; y < length; y++) {  
            s += arr[x] + arr[y];  
        }  
    }  
    return s;  
}
```

**Constant time, done $\text{length} * \text{length}$ times.
Call this C_7 .**

Another Example

```
int sum2(int* arr, int length) {  
    int s = 0, x, y;  
    for (x = 0; x < length; x++) {  
        for (y = 0; y < length; y++) {  
            s += arr[x] + arr[y];  
        }  
    }  
    return s;  
}
```

Constant time, done $\text{length} * \text{length}$ times.
Call this C_8 .

Another Example

```
int sum2(int* arr, int length) {  
    int s = 0, x, y;  
    for (x = 0; x < length; x++) {  
        for (y = 0; y < length; y++) {  
            s += arr[x] + arr[y];  
        }  
    }  
    return s;  
}
```

Constant time, done once. Call this C_9 .

Putting it Together

- We are left with the following formula:

$$c_1 + c_2 + (\text{length} * c_3) + (\text{length} * c_4) + (\text{length} * c_5) + (\text{length} * \text{length} * c_6) + (\text{length} * \text{length} * c_7) + (\text{length} * \text{length} * c_8) + c_9$$

Putting it Together

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$c_1 + c_2 + (\text{length} * c_3) + (\text{length} * c_4) + (\text{length} * c_5) + (\text{length} * \text{length} * c_6) + (\text{length} * \text{length} * c_7) + (\text{length} * \text{length} * c_8) + c_9$$

Putting it Together

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$\begin{aligned} &1 + 1 + (\text{length} * 1) + (\text{length} * 1) + \\ &(\text{length} * 1) + (\text{length} * \text{length} * 1) + \\ &\quad (\text{length} * \text{length} * 1) + \\ &\quad (\text{length} * \text{length} * 1) + 1 \end{aligned}$$

Putting it Together

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$3 + \text{length} + \text{length} + \text{length} + (\text{length} * \text{length}) + (\text{length} * \text{length}) + (\text{length} * \text{length})$$

Putting it Together

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$3 + 3(\text{length}) + 3(\text{length} * \text{length})$$

Putting it Together

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

`1 + length + (length * length)`

Putting it Together

- The specific values of constants are unimportant as long as they are positive
- We can replace all these with the value 1 as far as Big O notation is concerned

$$1 + \text{length} + \text{length}^2$$

Putting it Together

- With sums, we always choose the larger sum
- A variable is always larger than a constant

$$1 + \text{length} + \text{length}^2$$

Putting it Together

- With sums, we always choose the larger sum
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$$\text{length} + \text{length}^2$$

Putting it Together

- With sums, we always choose the larger sum
- A variable is always larger than a constant

length²

Putting it Together

- Observe that `length` is really N , the input size
- For this example, we are done

`length2`

Putting it Together

- Observe that `length` is really N , the input size
- For this example, we are done

$$O(N^2)$$

Big O Heuristics

- A non-loop is often $O(1)$
- A single loop is often $O(N)$
- A singly nested loop is often $O(N^2)$
- Not always true though - we will see exceptions later in this class
- Determining time complexity can be quite difficult in general